

Robust and Data-Driven Markov Decision Processes

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Markov decision process

Tuple $(\mathcal{S}, \mathcal{A}, q, p, r, \lambda)$ where

- $\mathcal{S} = \{1, \dots, S\}$ is the (finite) **state space**;
- $\mathcal{A} = \{1, \dots, A\}$ is the (finite) **action space**;
- $q = (q_1, \dots, q_S) \in \Delta(\mathcal{S})$ is the **initial state distribution**;
- $p : \mathcal{S} \times \mathcal{A} \rightarrow \Delta(\mathcal{S})$ is the **transition kernel** with elements $p(s' | s, a)$;
- $r : \mathcal{S} \times \mathcal{A} \rightarrow \mathbb{R}$ are the **expected one-step rewards**;
- $\lambda \in (0, 1)$ is the **discount factor**.

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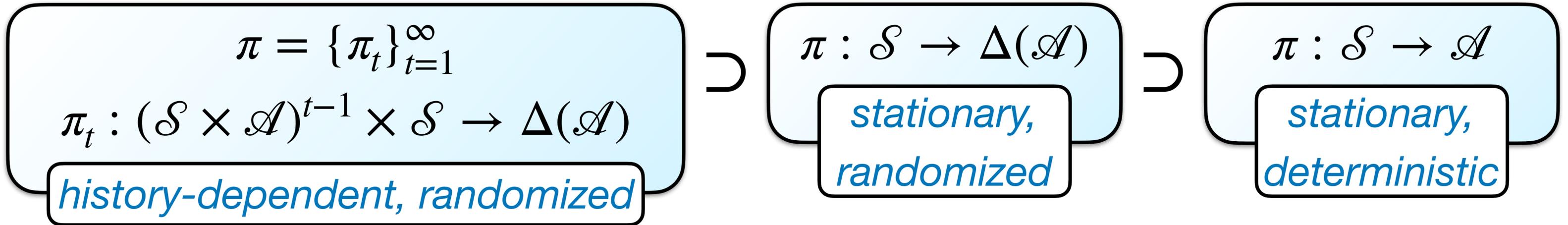
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- $r : \mathcal{S} \times \mathcal{A} \rightarrow \mathbb{R}$ is the reward function;
- $\lambda \in (0, 1)$ is the discount factor.

Objective

find **policy** π that maximizes the expected total discounted rewards:

$$\underset{\pi \in \Pi}{\text{maximize}} \quad \mathbb{E}_p \left[\sum_{t=1}^{\infty} \lambda^{t-1} \cdot r(s_t, \pi_t[s_t]) \right]$$

- **Stationary deterministic policies are optimal:**



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$$v^\pi(s) = \mathbb{E}_p \left[\sum_{t=1}^{\infty} \lambda^{t-1} \cdot r(s_t, \pi[s_t]) \mid s_1 = s \right]$$

- **Stationary deterministic policies are optimal.**
- **Discounted rewards of a fixed policy satisfy linear equations:**

$$v^\pi(s) = r(s, \pi[s]) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, \pi[s]) \cdot v^\pi(s')$$

- **Stationary deterministic policies are optimal.**
- **Discounted rewards of a fixed policy satisfy linear equations.**
- **Discounted rewards of an optimal policy satisfy nonlinear equations:**

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

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- **v^* -greedy policy is optimal:**

$$\pi^*(s) \in \arg \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

- **Value iteration:**

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

- Value iteration:

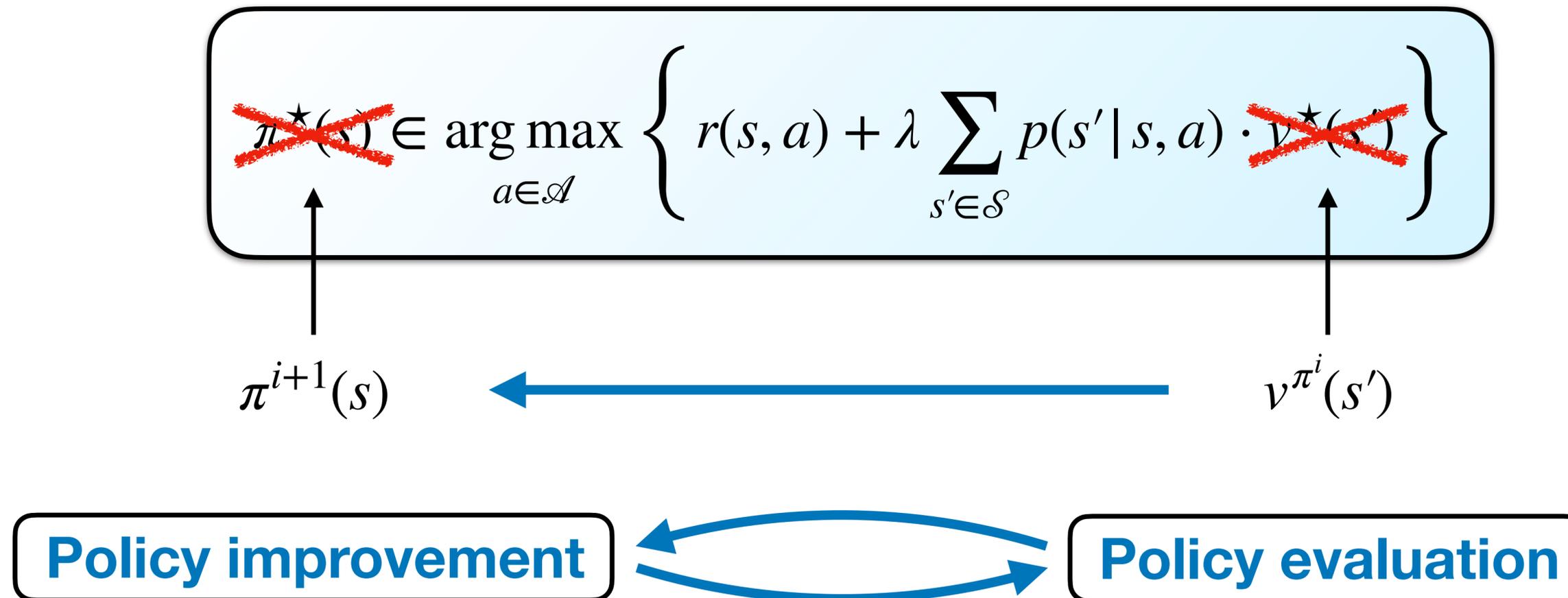
$$v^{i+1}(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^i(s') \right\}$$

Starting from any $v^0 \in \mathbb{R}^{\mathcal{S}}$, converges at *linear rate* to v^{\star} .

- **Value iteration.**
- **(Modified) Policy iteration:**

$$\pi^*(s) \in \arg \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

- Value iteration.
- **(Modified) Policy iteration:**



Under suitable conditions, *converges at superlinear rate* to v^* .
Converges to (ϵ -)optimal policy in finitely many iterations.

- **Value iteration.**
- **(Modified) Policy iteration.**
- **Linear programming:**

$$\begin{array}{ll} \text{minimize} & \sum_{s \in \mathcal{S}} q(s) \cdot v(s) \\ \text{subject to} & v(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \quad \forall s \in \mathcal{S} \end{array}$$

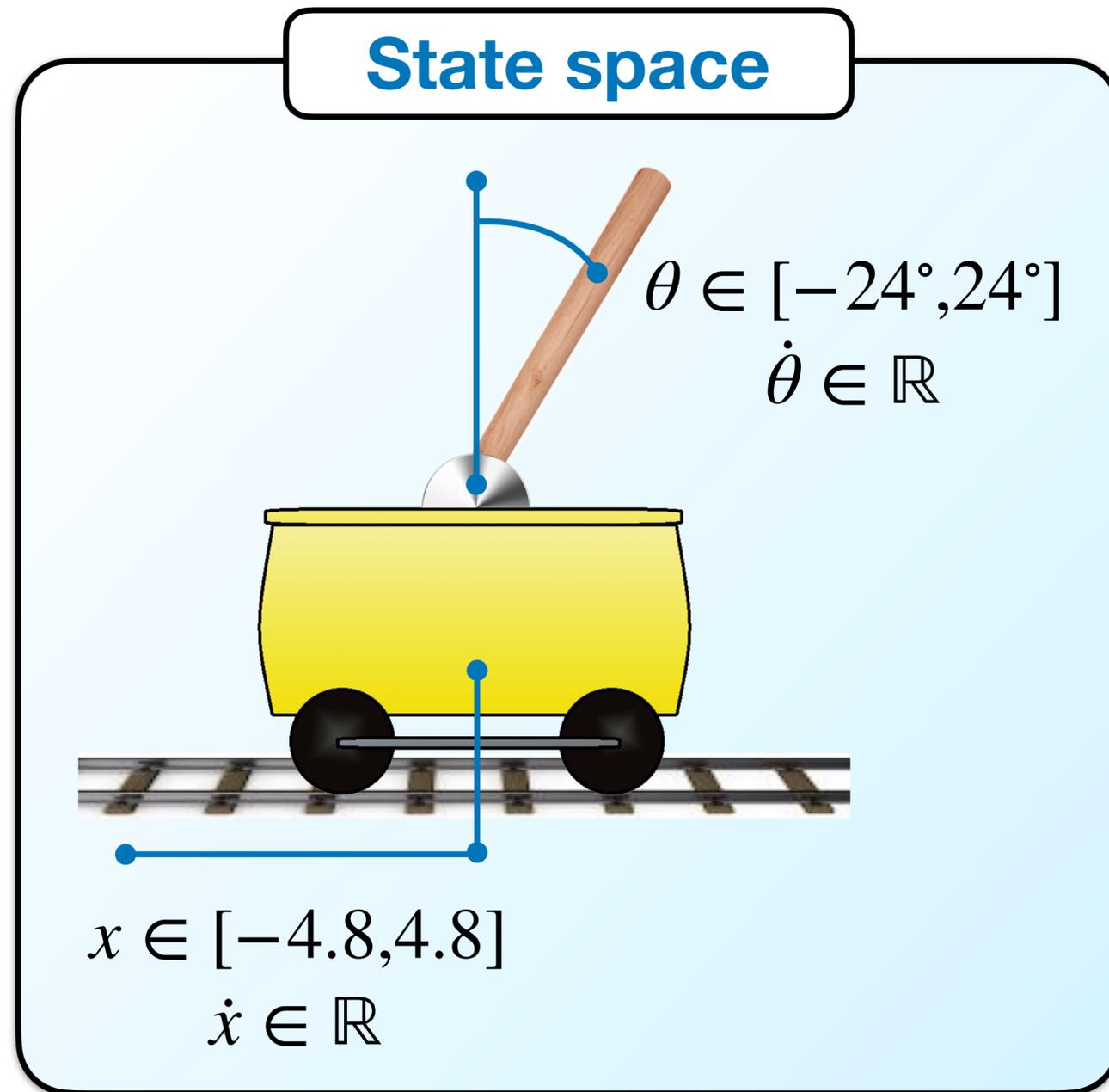
- Value iteration.
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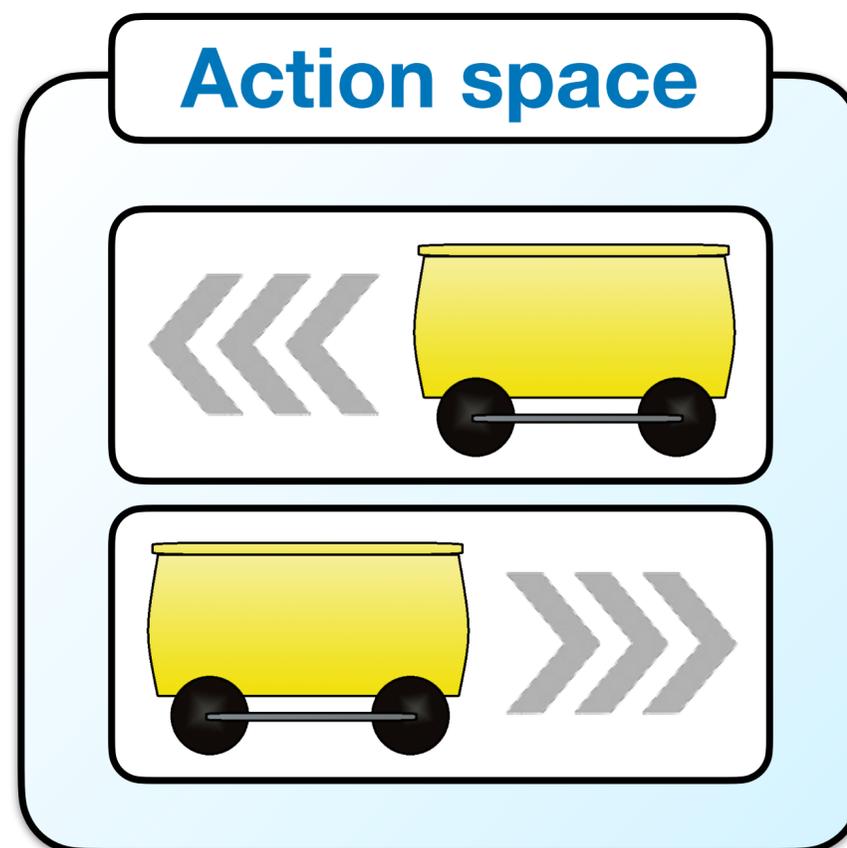
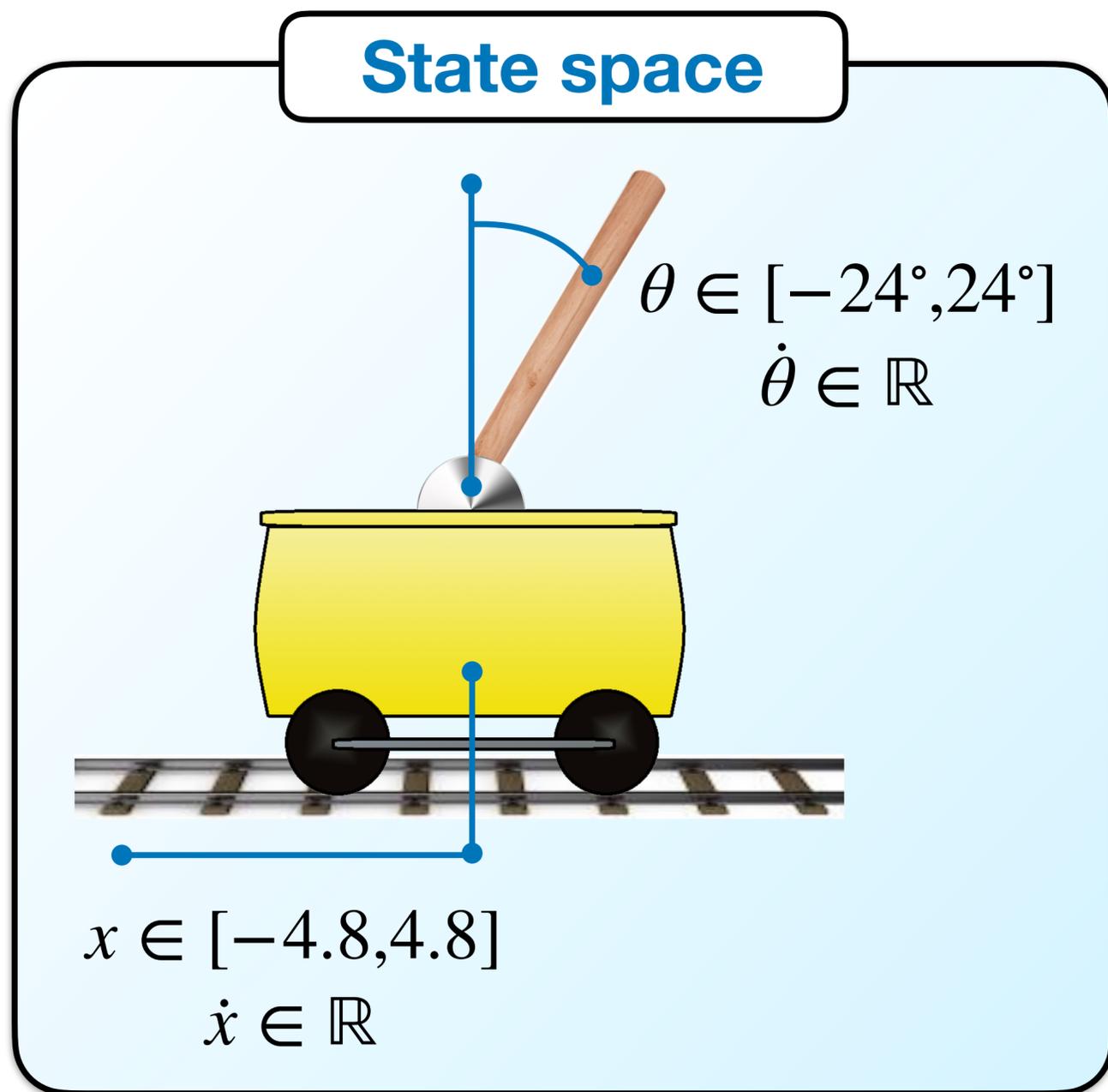
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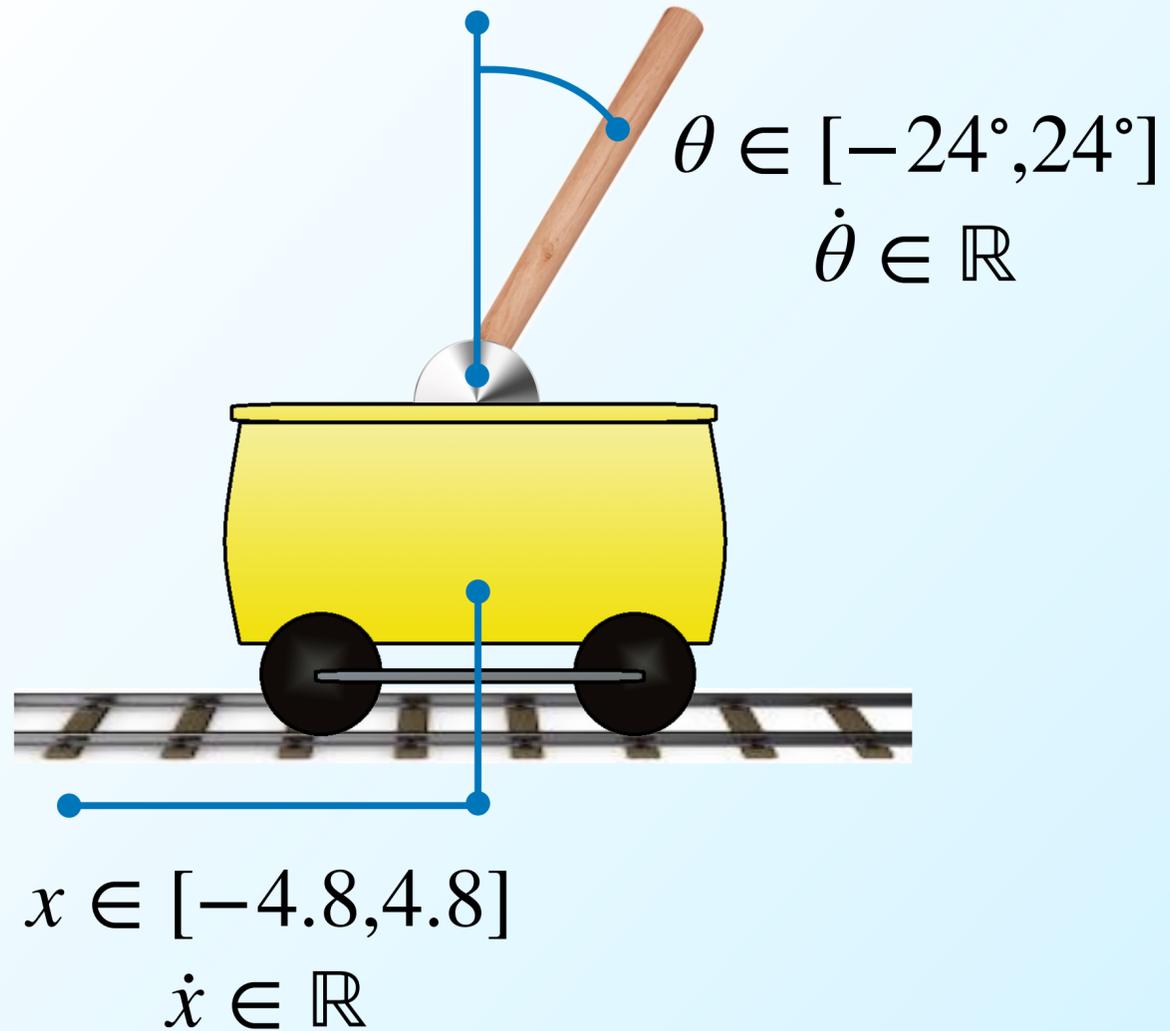
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Optimally solved in polynomial time with standard solvers.

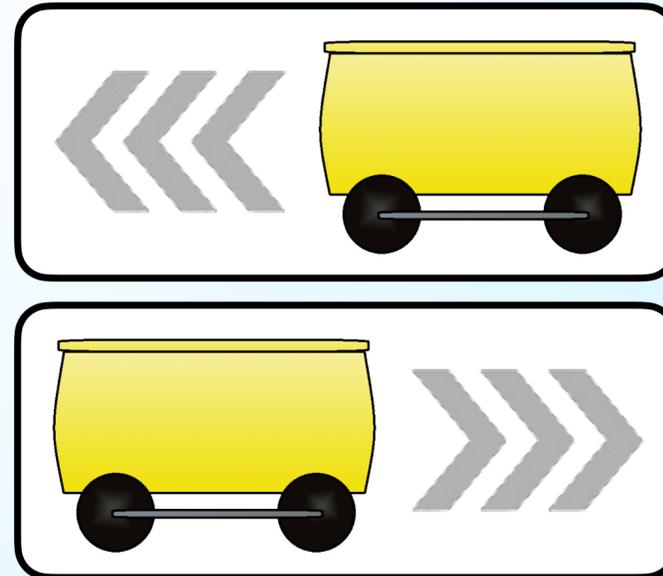




State space



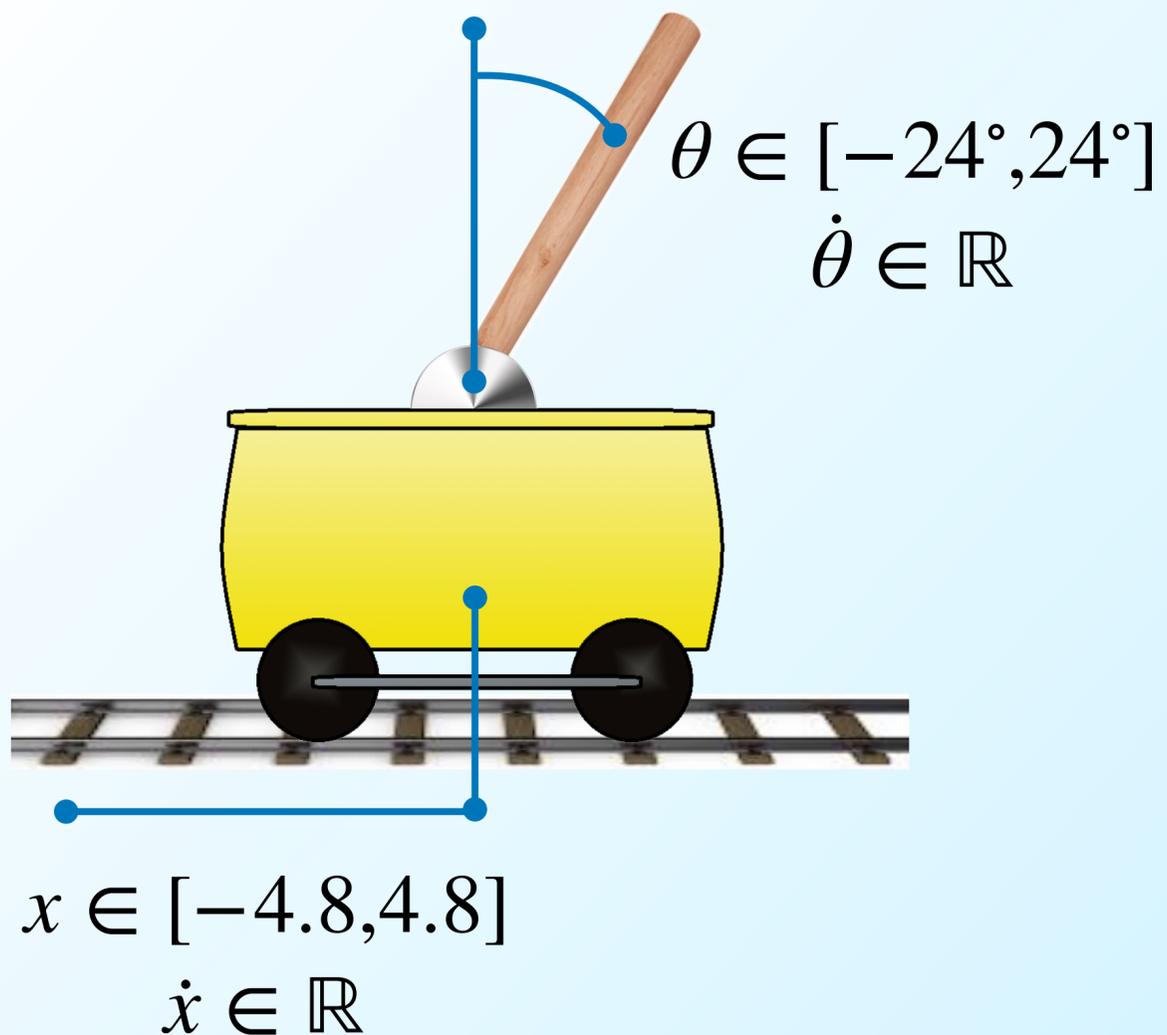
Action space



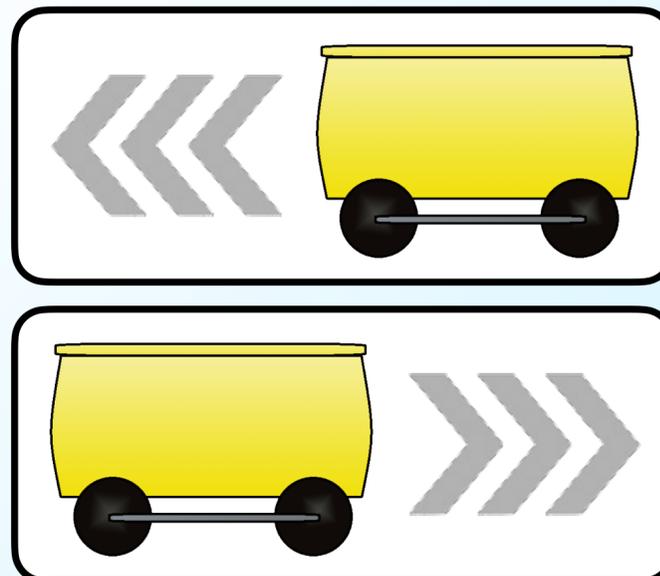
Initial state

$$x, \dot{x}, \theta, \dot{\theta} \sim \mathcal{U}[-0.05, 0.05]$$

State space



Action space



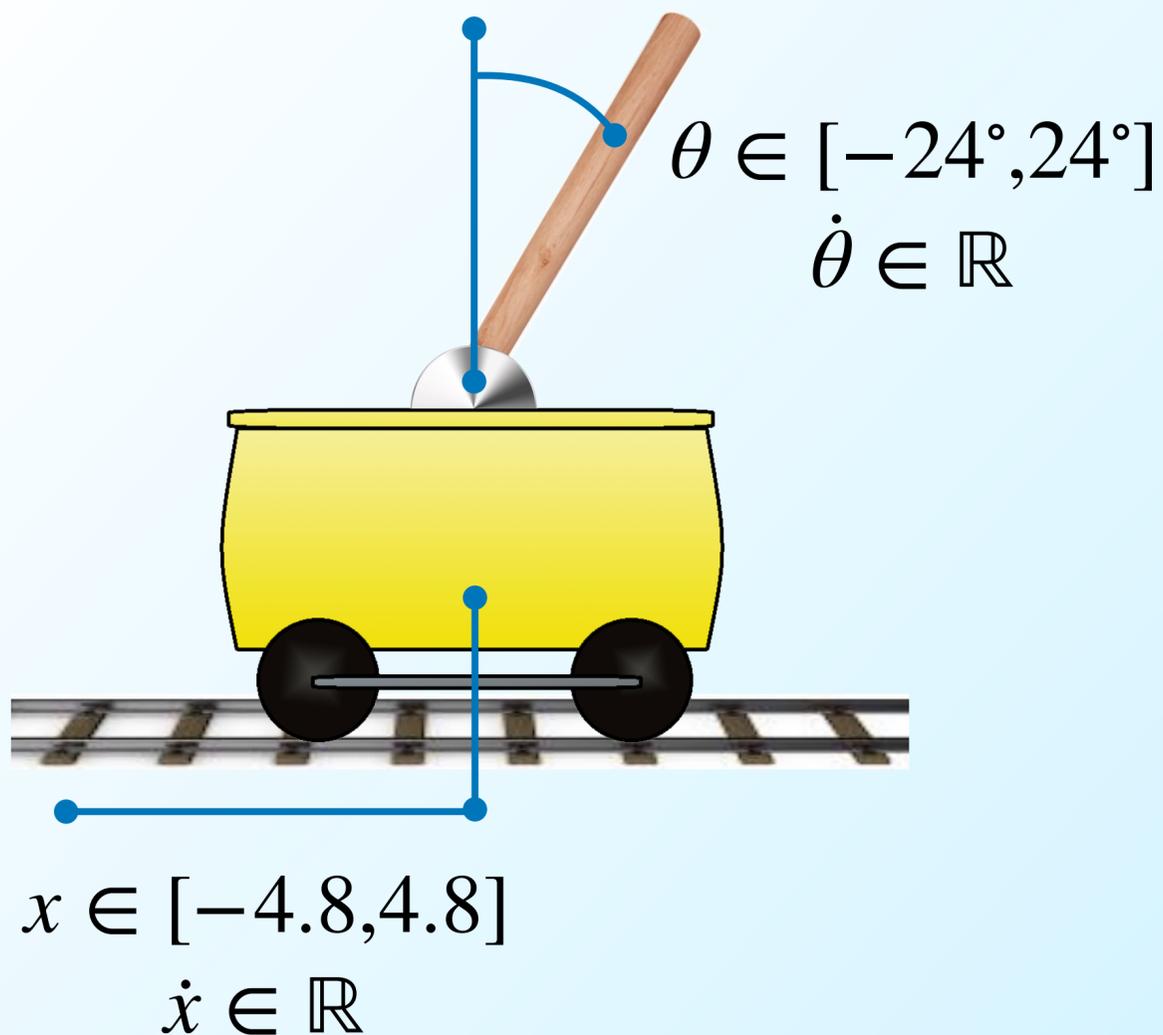
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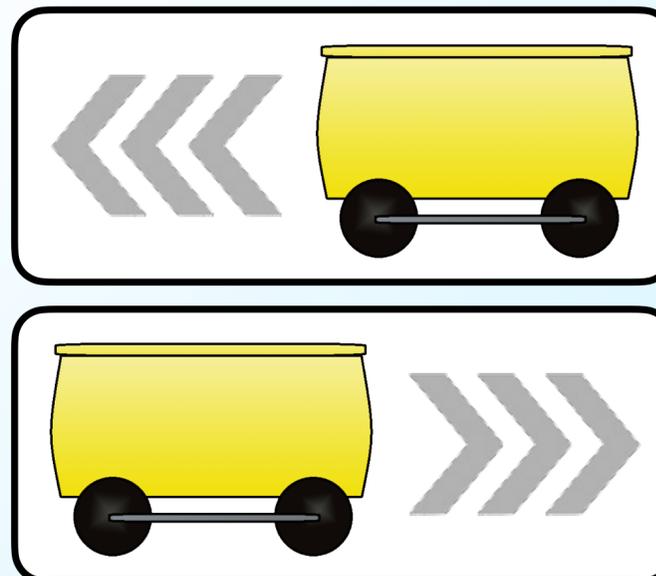
Transitions

- deterministic via laws of mechanics
- terminate if
$$x \notin [-2.4, 2.4]$$
or $\theta \notin [-12^\circ, 12^\circ]$

State space



Action space



Initial state

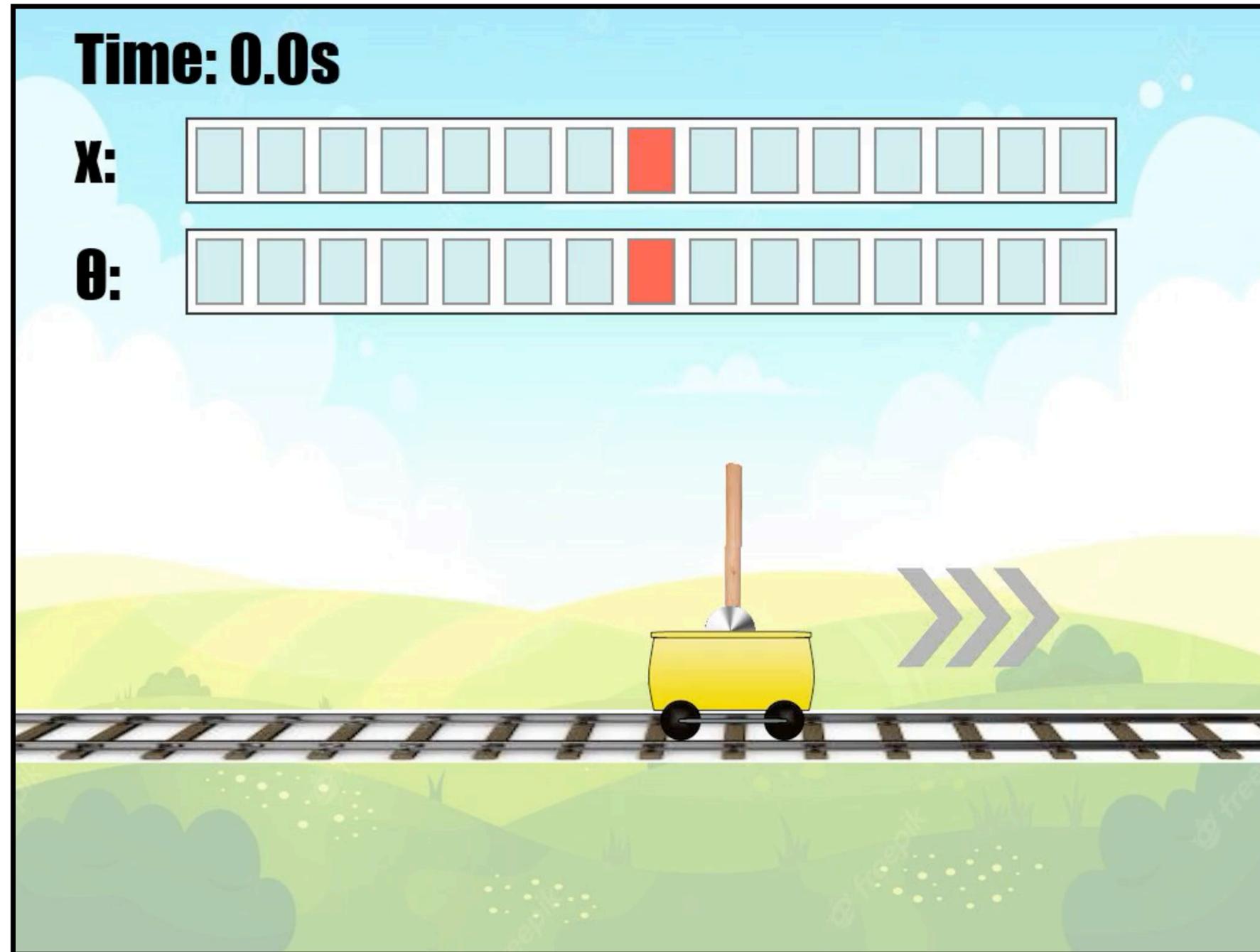
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Rewards

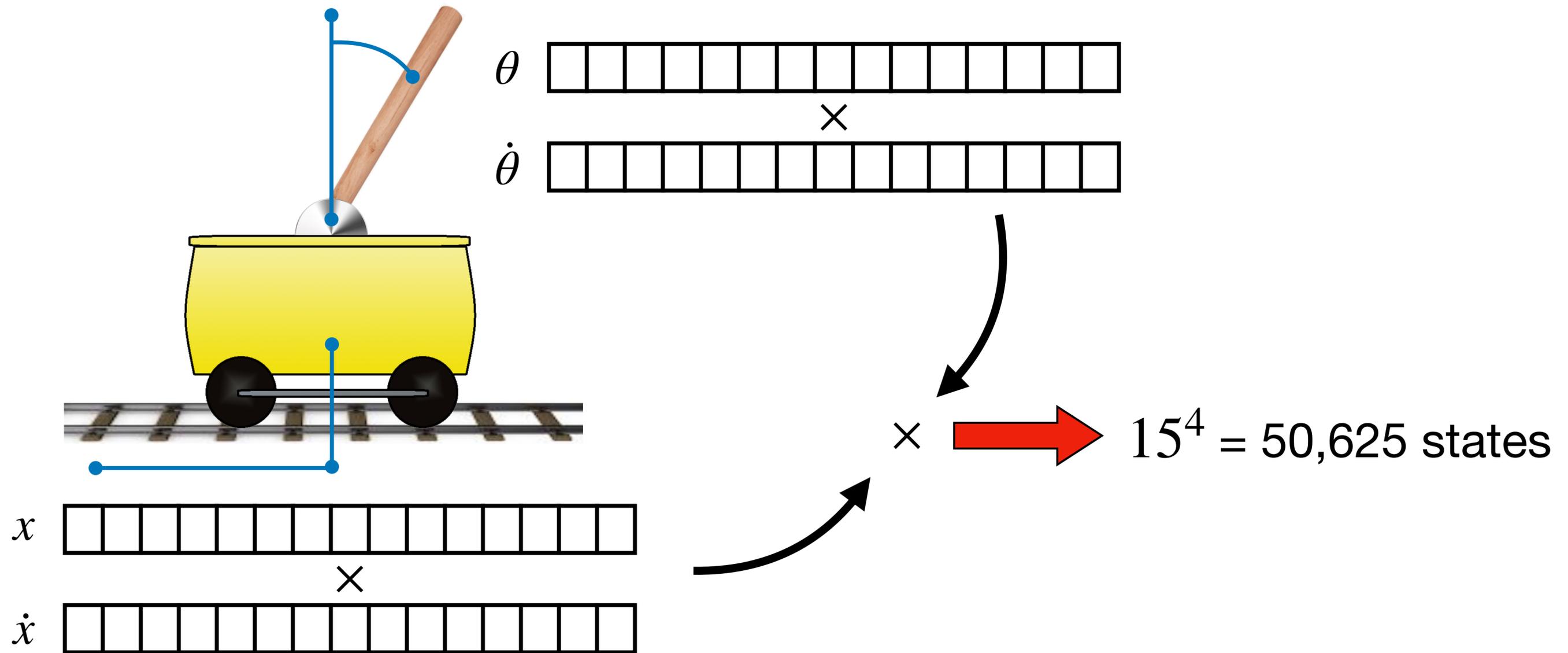
+1/non-terminated time step



Ambiguity and Robust MDPs

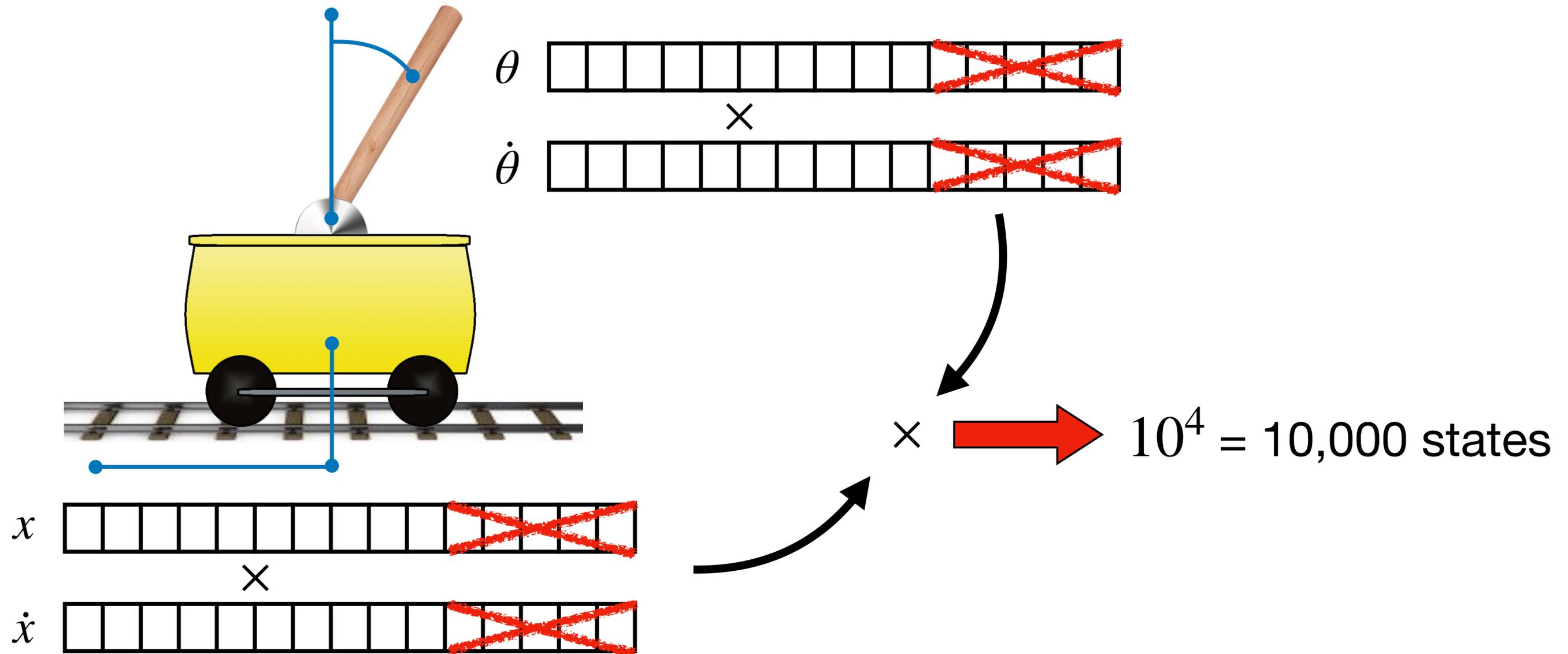
Two common sources of **ambiguity**:

- **Modelling errors**: 32.67 secs/run



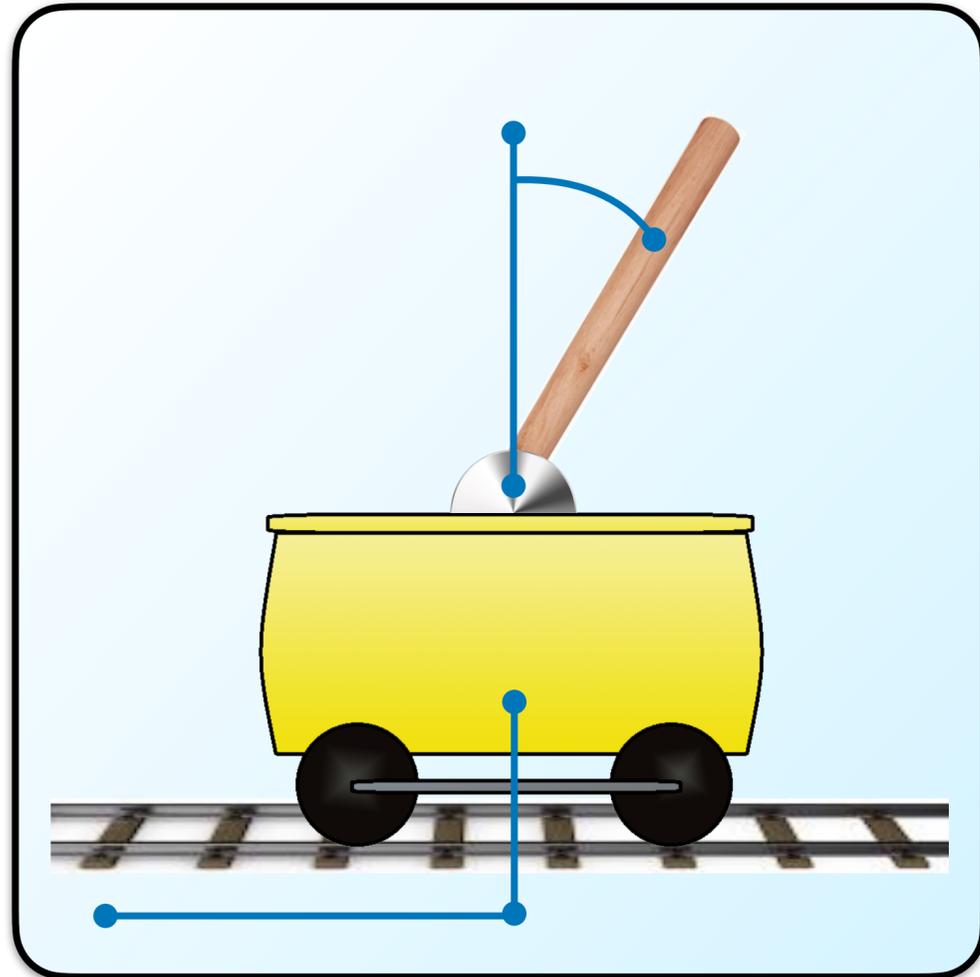
Two common sources of **ambiguity**:

- **Modelling errors**: 32.67 secs/run \rightarrow 2.45 secs/run



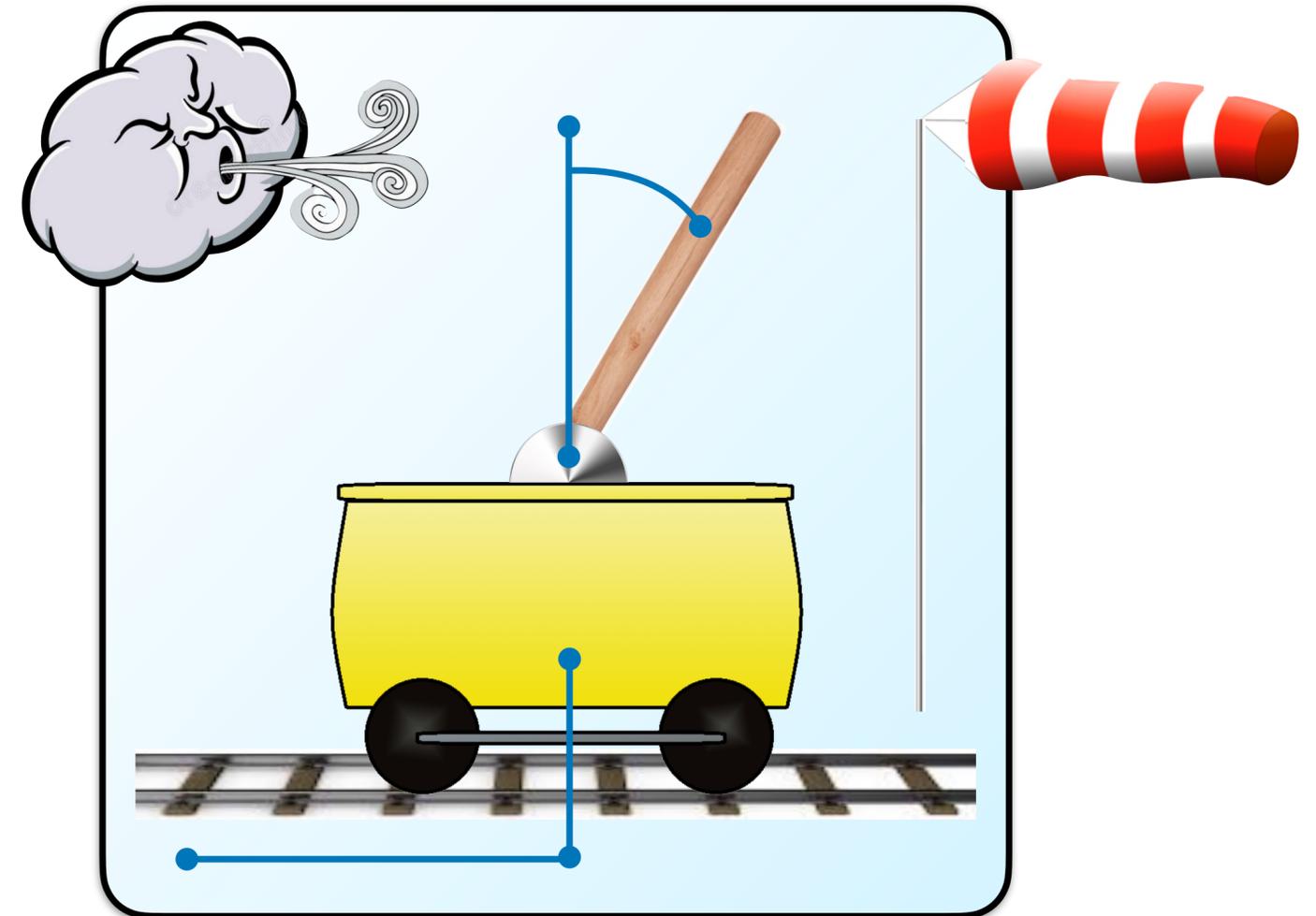
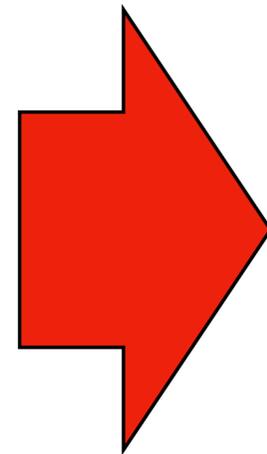
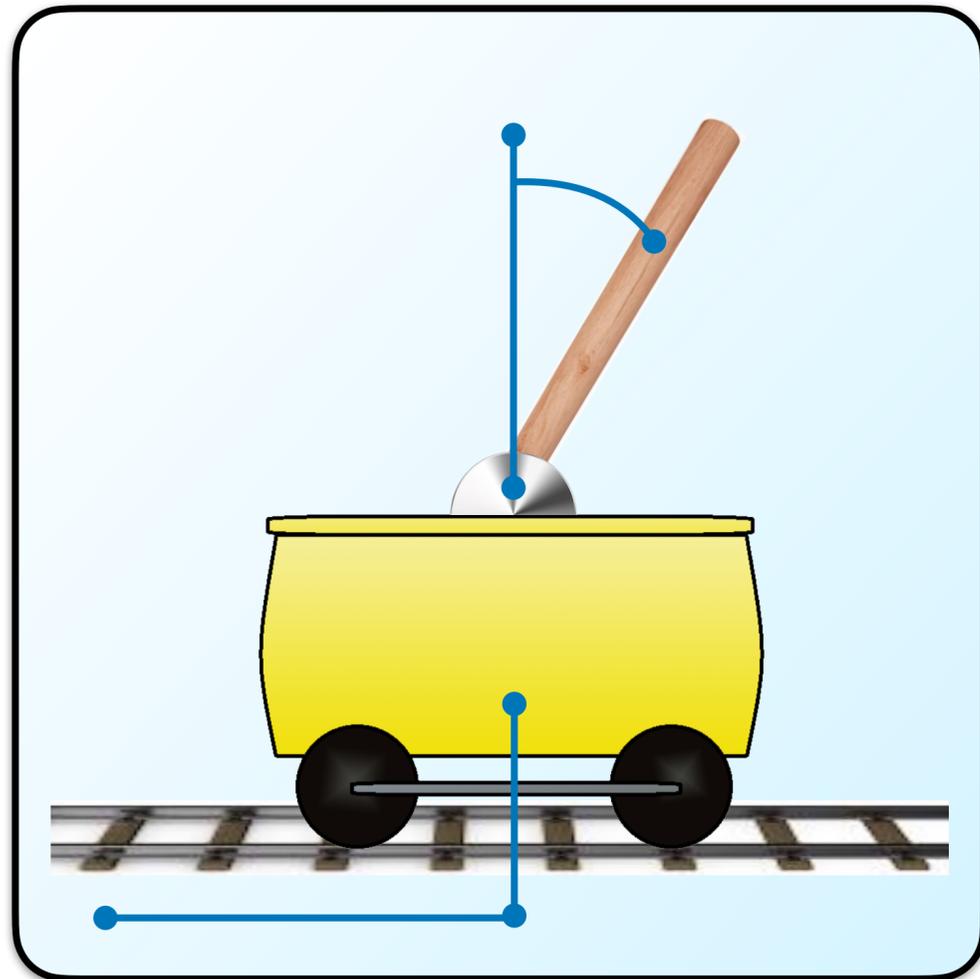
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- **Estimation errors**: 32.67 secs/run



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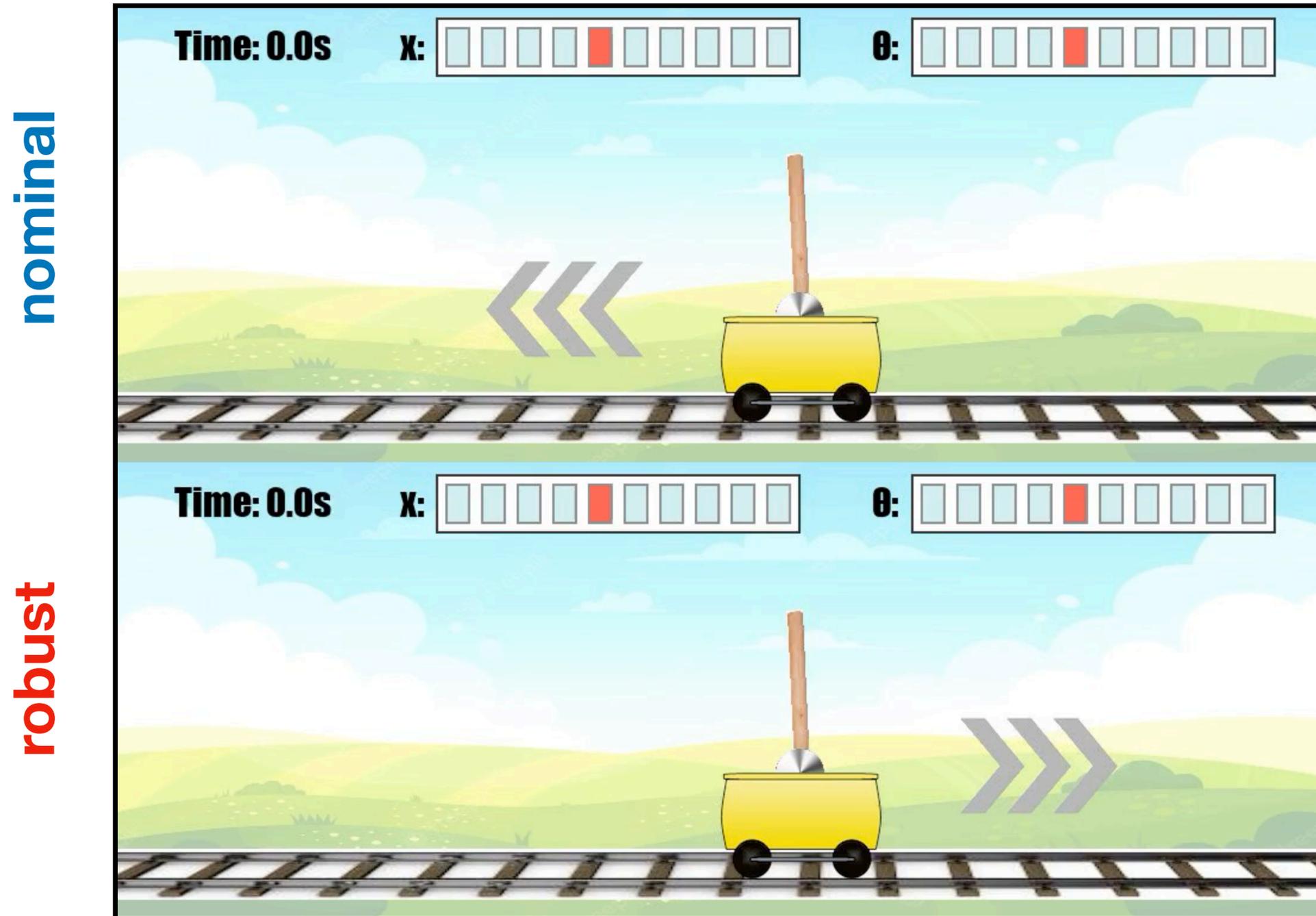
Impact of **ambiguity** can be alleviated via **robust optimization**:

Robust MDP

$$\underset{\pi \in \Pi}{\text{maximize}} \quad \underset{p \in \mathcal{P}}{\text{inf}} \quad \mathbb{E}_p \left[\sum_{t=1}^{\infty} \lambda^{t-1} \cdot r(s_t, \pi[s_t]) \right]$$

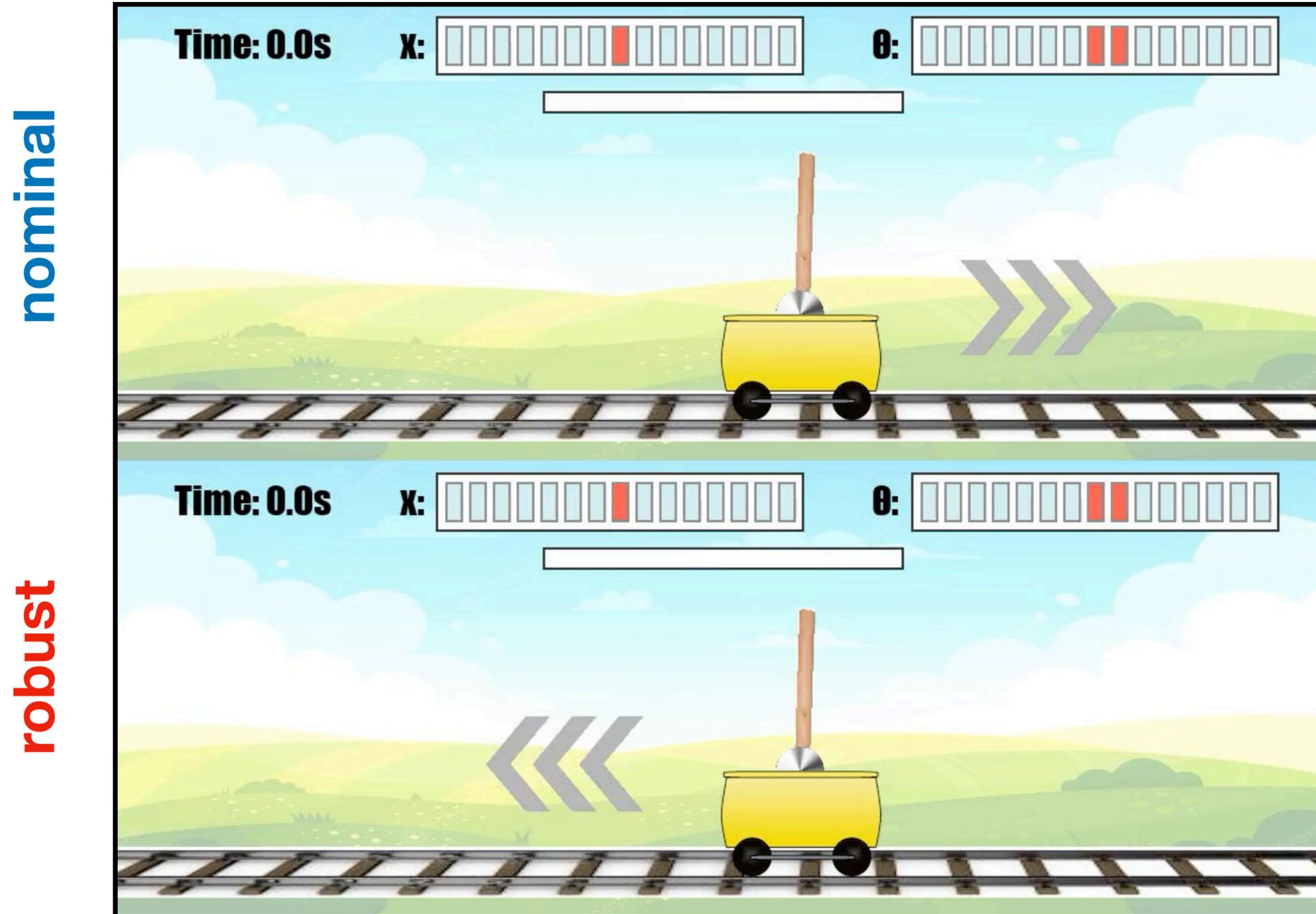
\searrow **ambiguity set**

Robust MDPs admit interpretation as regularized MDPs!



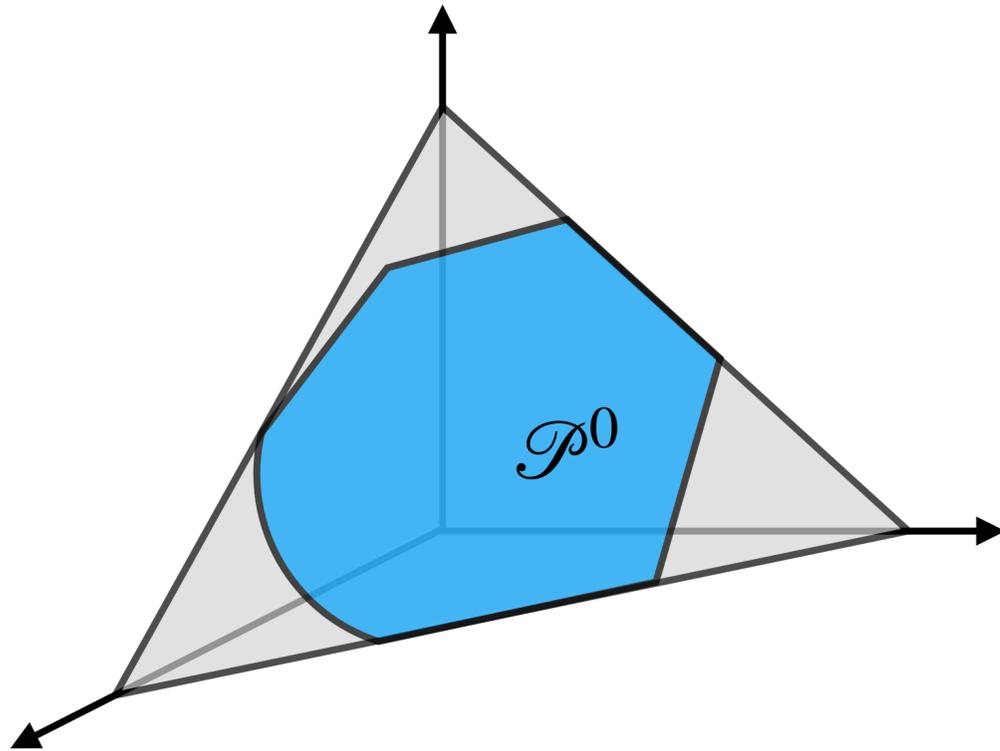
Modelling errors: 32.67 secs/run → 2.45 secs/run → 15.77 secs/run

Ambiguity: Estimation Errors



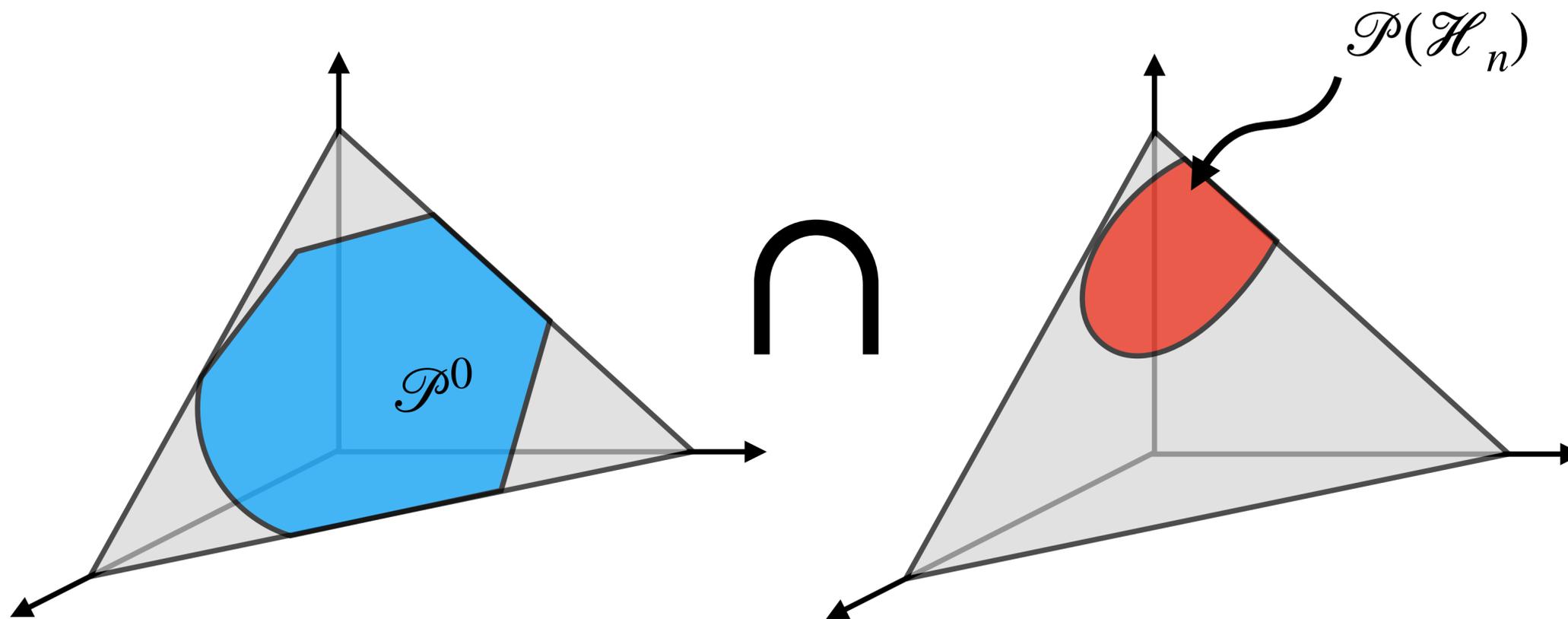
Estimation errors: 32.67 secs/run → 4.68 secs/run → 15.76 secs/run

Structural ambiguity set



Structural ambiguity set

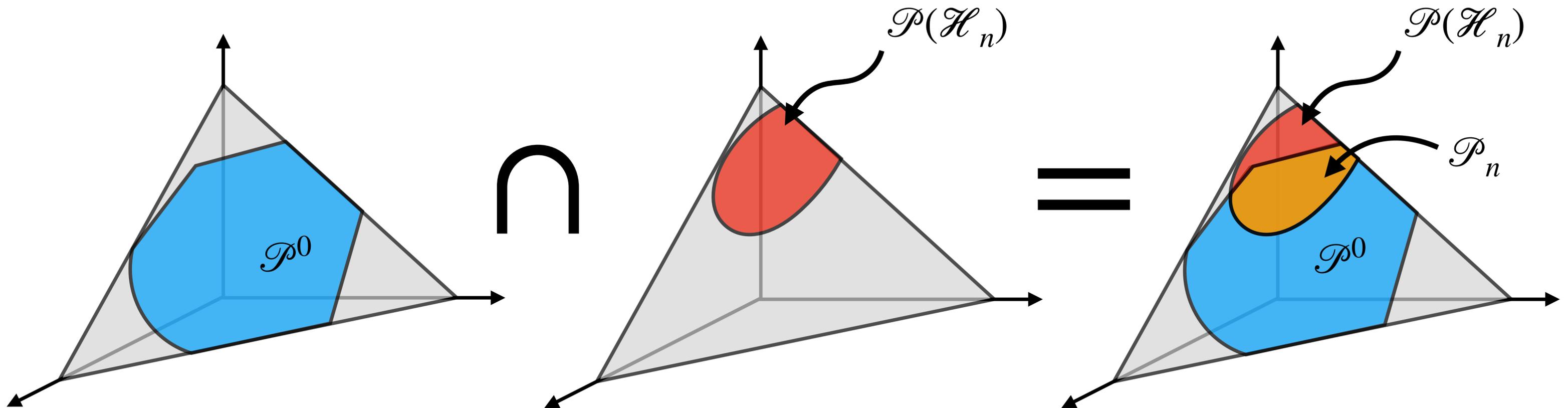
Historical sample



Structural ambiguity set

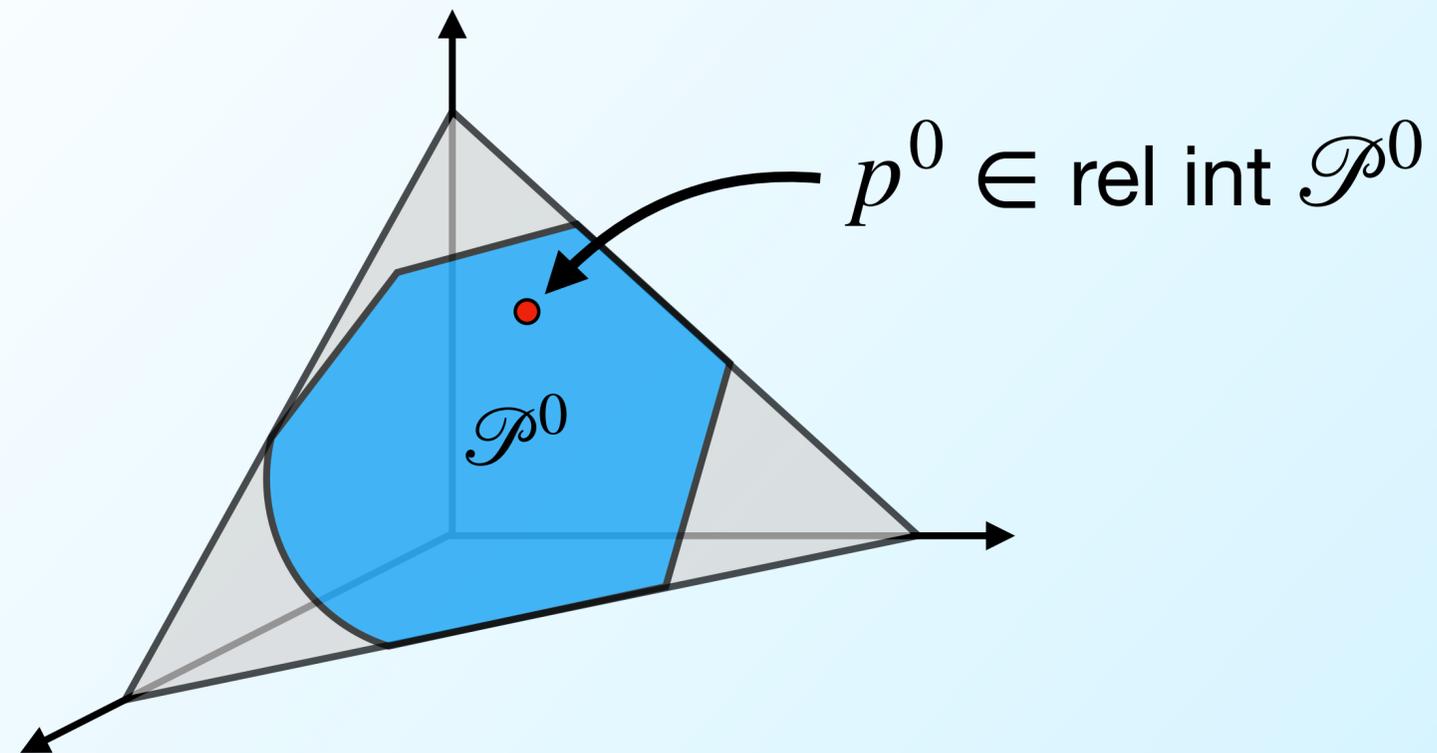
Historical sample

Out-of-sample guarantee



Structural ambiguity set

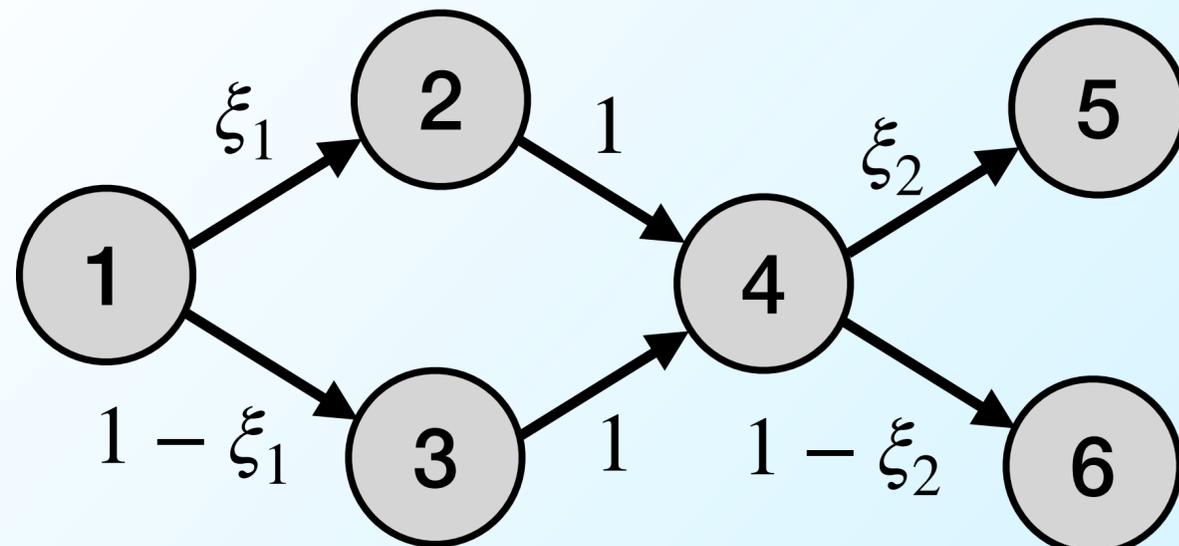
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Structural ambiguity set

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Possible transitions



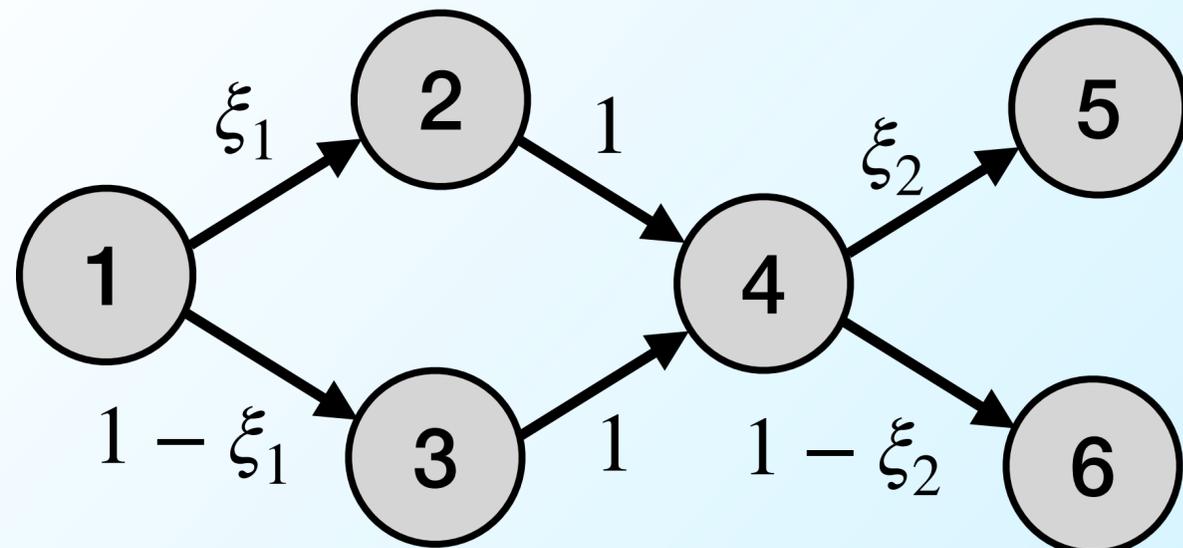
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Structural ambiguity set

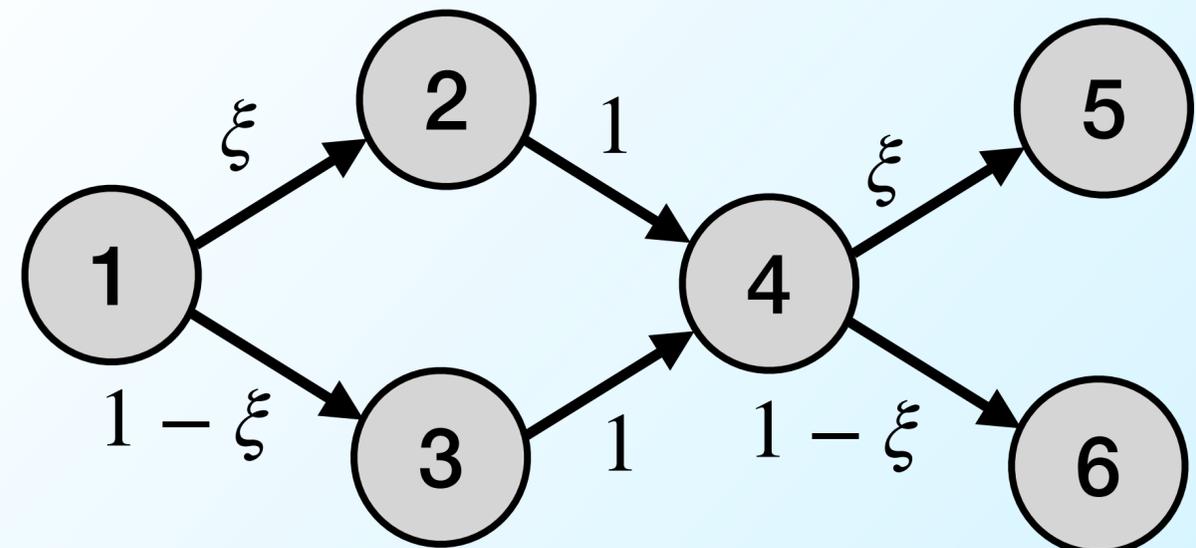
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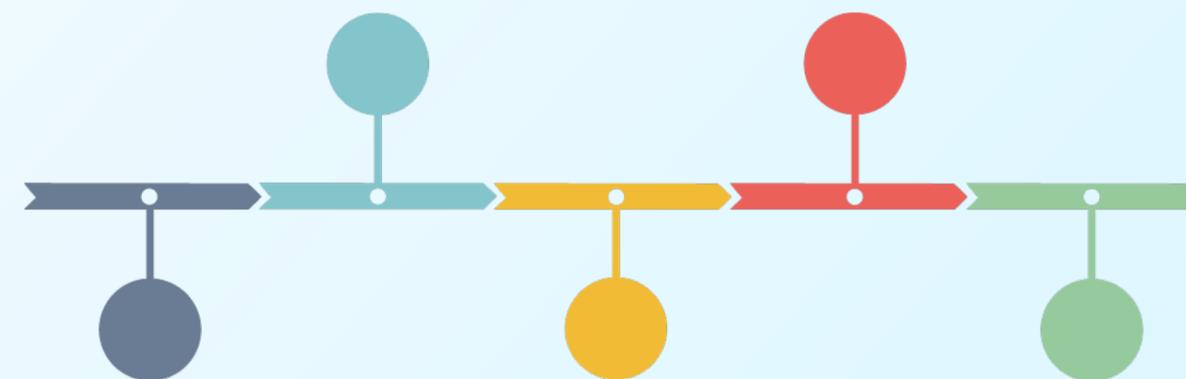
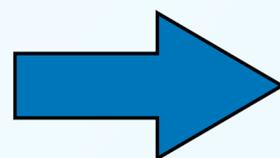
Possible transitions



Equal probabilities



Historical sample

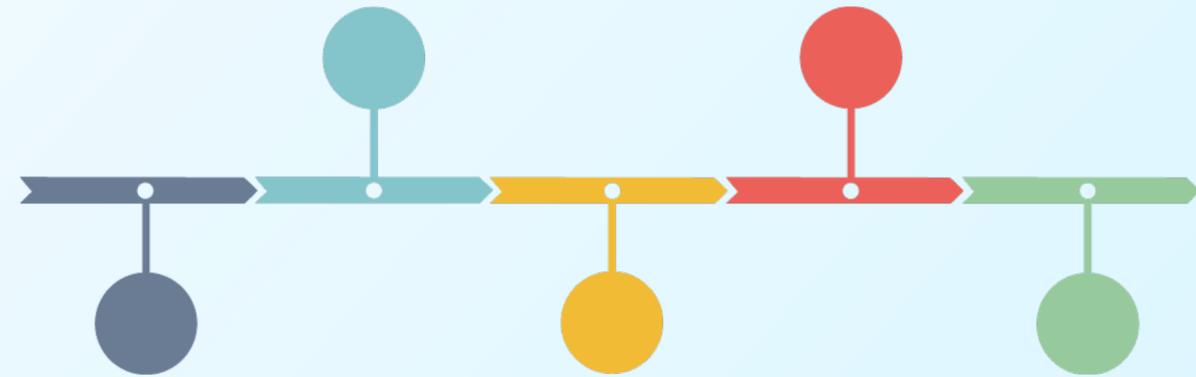
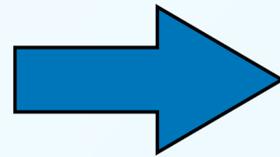


historical policy π^0
(stationary, randomized)

state-action history

$$\mathcal{H}_n = (s_1, a_1, \dots, s_n, a_n) \in (\mathcal{S} \times \mathcal{A})^n$$

Historical sample



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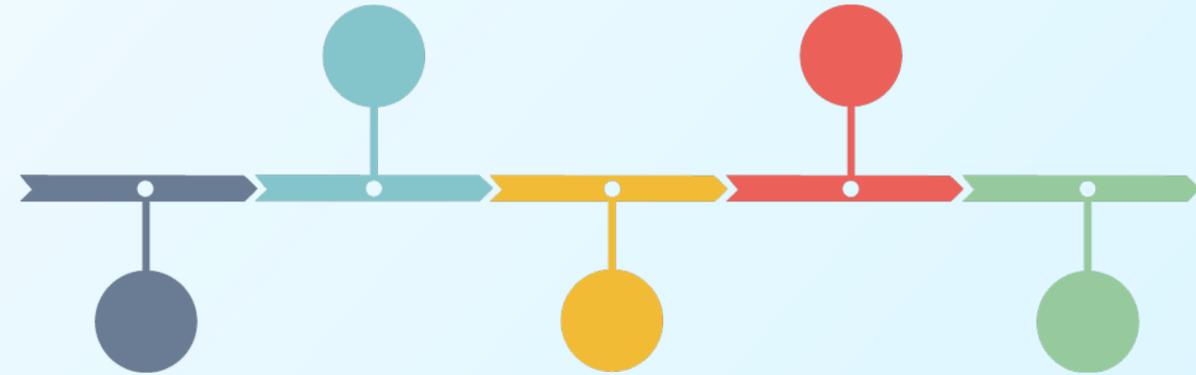
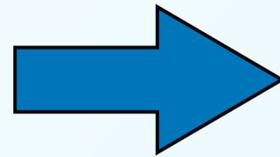
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Likelihood, given history

$$\mathcal{L}_n(p) = q(s_1) \cdot \pi^0(a_n | s_n) \cdot \prod_{t=1}^{n-1} [\pi^0(a_t | s_t) \cdot p(s_{t+1} | s_t, a_t)]$$

Historical sample

$$\mathcal{P}(\mathcal{H}_n) = \{p : \log \mathcal{L}_n(p) \geq \log \mathcal{L}_n(p^*) - \delta\}$$



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Theorem

Assumption: Historical policy π^0 visits every $s \in \mathcal{S}$ infinitely often as $n \rightarrow \infty$

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$\mathcal{P}_n = \mathcal{P}^0 \cap \mathcal{P}(\mathcal{H}_n)$ with $\delta = (1 - \beta)$ -quantile of χ^2 -distribution with κ degrees of freedom

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$$\text{plim}_{n \rightarrow \infty} \left[\sqrt{n} \cdot d^{\text{H}}(\mathcal{P}_n, \{p^0\}) \right] = 0$$

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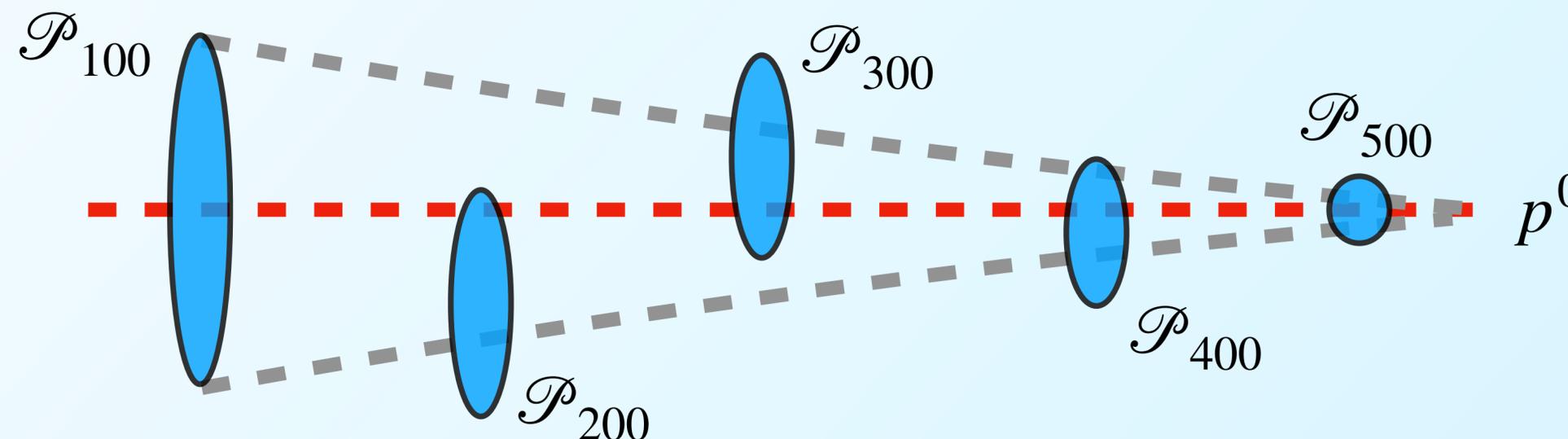
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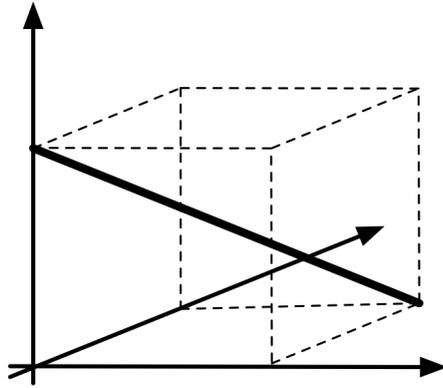
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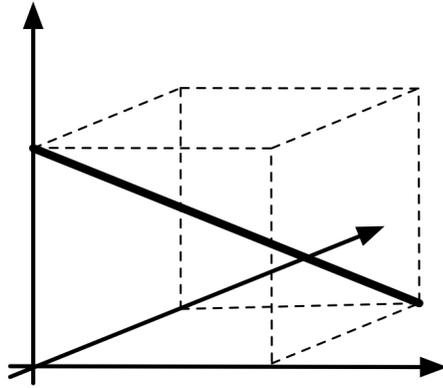
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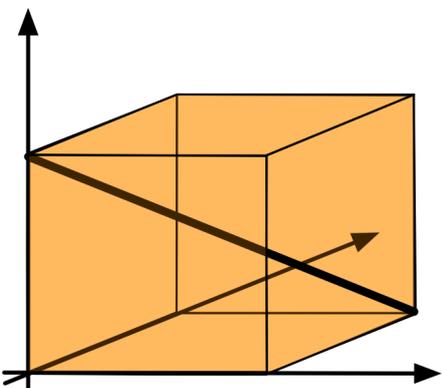
General (non-rectangular) ambiguity sets

- 👎 Optimal policy can be **randomized** & **history-dependent**
- 👎 Bellman optimality principle **violated**; **NP-hard**



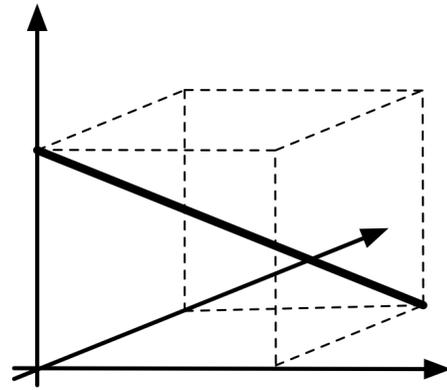
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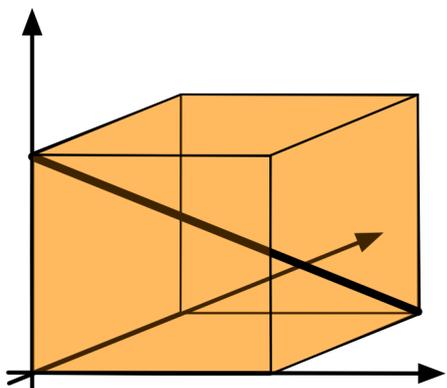
(s,a)-rectangular ambiguity sets

$$\mathcal{P} = \prod_{(s,a) \in \mathcal{S} \times \mathcal{A}} \mathcal{P}_{s,a} \quad \text{with} \quad \mathcal{P}_{s,a} \subseteq \Delta(\mathcal{S})$$



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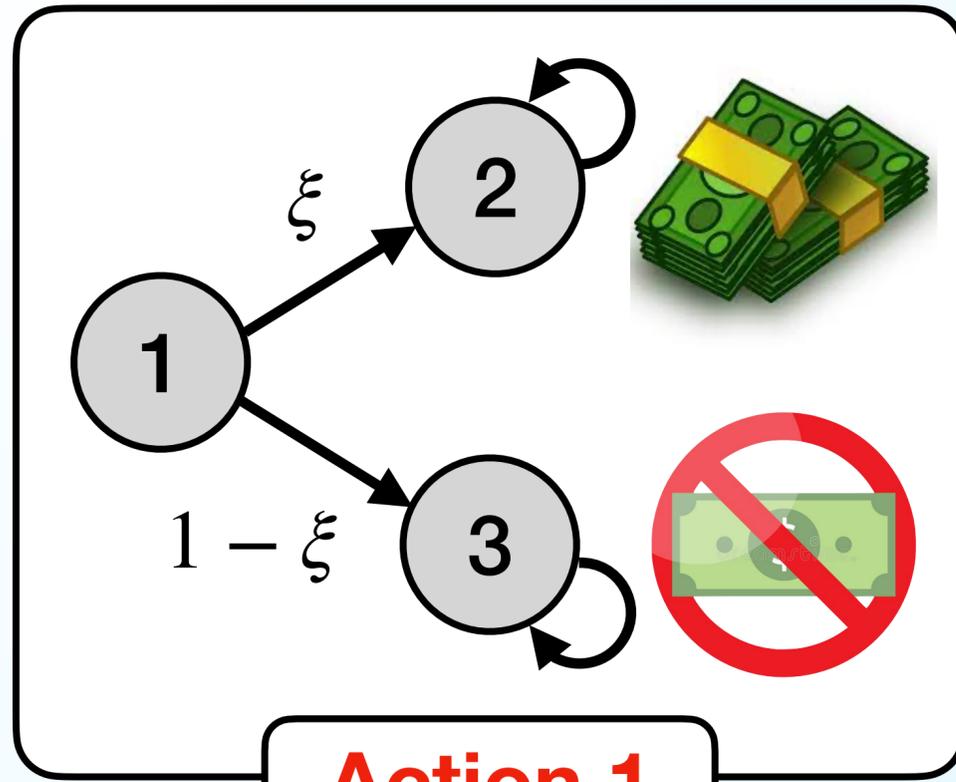


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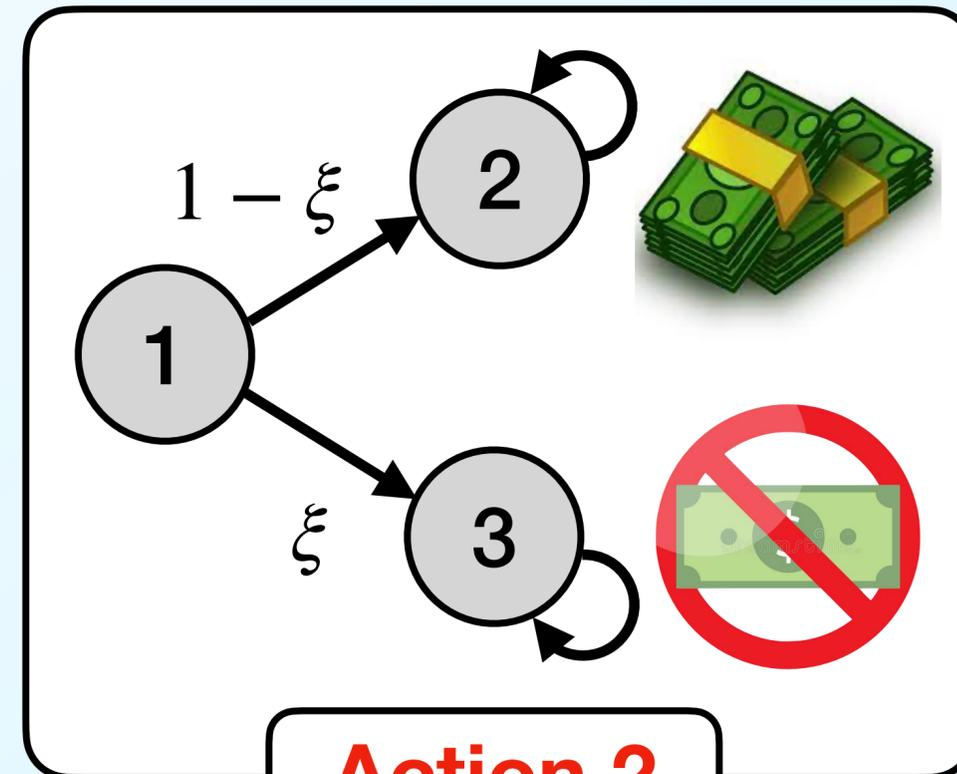
- 👍 Optimal policy **stationary** and **deterministic**
- 👍 Bellman optimality principle **holds**

General (non-rectangular) ambiguity sets

Example



Action 1



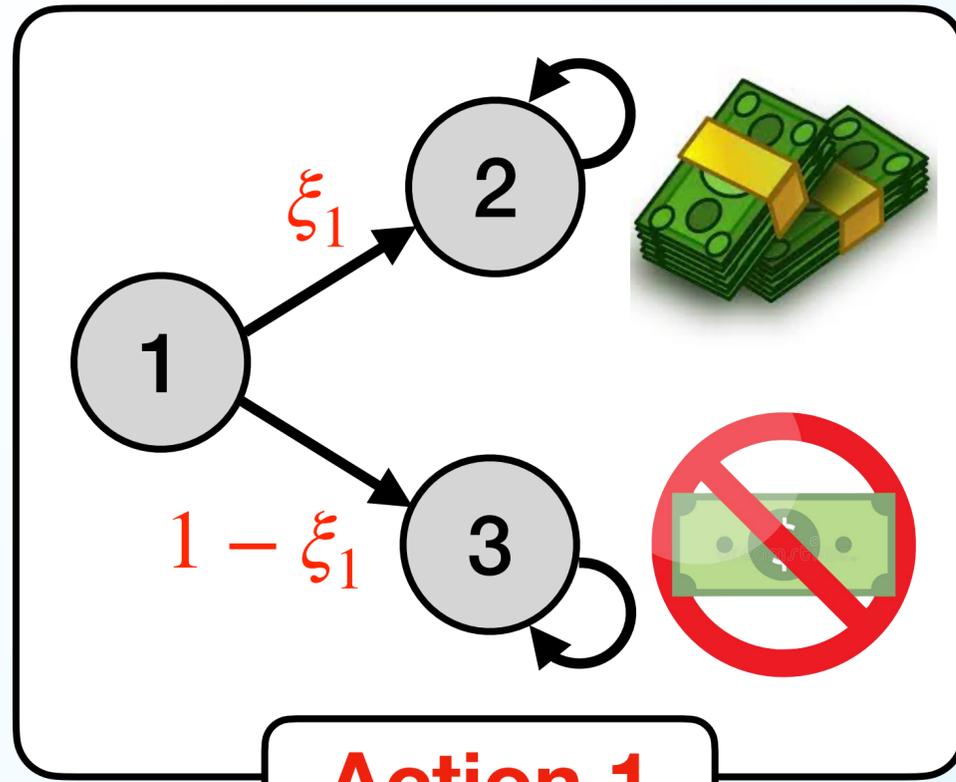
Action 2

for some unknown $\xi \in [0,1]$

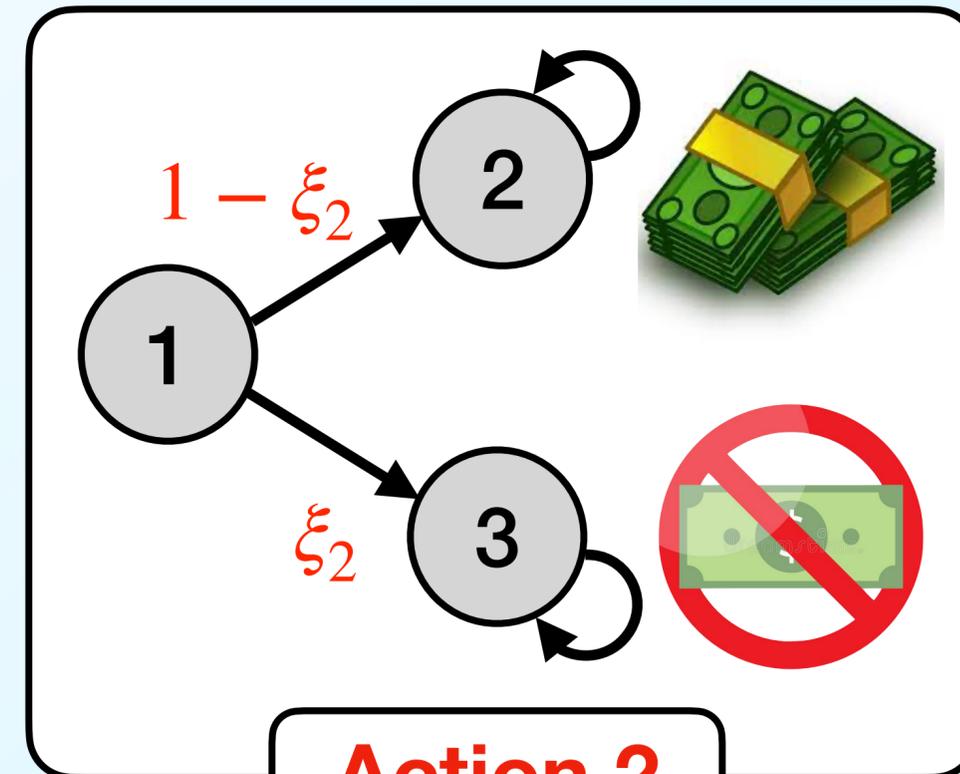
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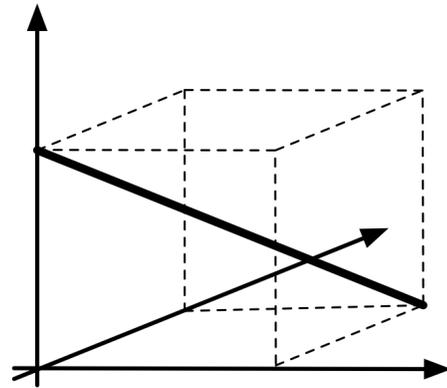


Action 2

for some unknown $\xi_1, \xi_2 \in [0,1]$

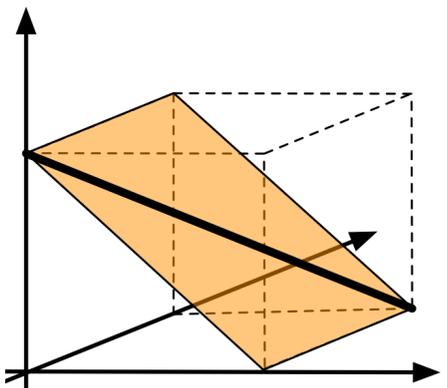
👊 Bellman optimality principle holds

endent



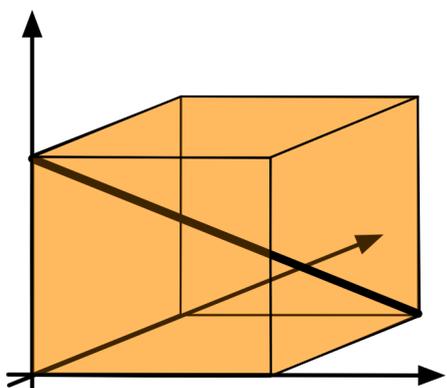
General (non-rectangular) ambiguity sets

- 👎 Optimal policy can be **randomized** & **history-dependent**
- 👎 Bellman optimality principle **violated**; **NP-hard**



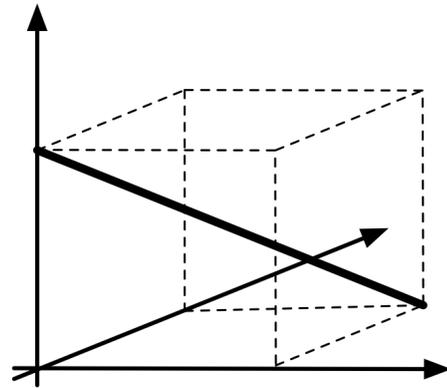
s-rectangular ambiguity sets

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s \subseteq [\Delta(\mathcal{S})]^A$$



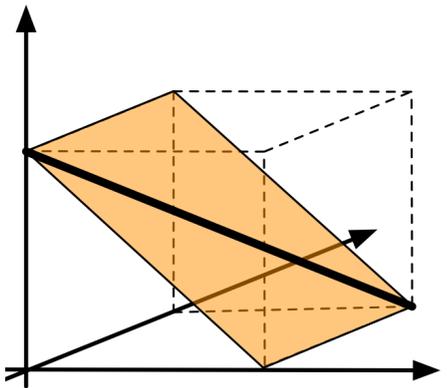
(s,a)-rectangular ambiguity sets

- 👍 Optimal policy **stationary** and **deterministic**
- 👍 Bellman optimality principle **holds**



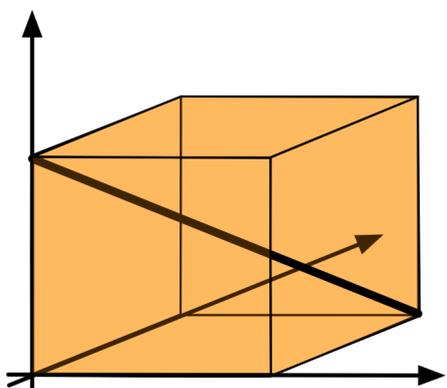
General (non-rectangular) ambiguity sets

- 👎 Optimal policy can be **randomized** & **history-dependent**
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s-rectangular ambiguity sets

- 👍 Optimal policy **stationary** but can be **randomized**
- 👍 Bellman optimality principle **holds**

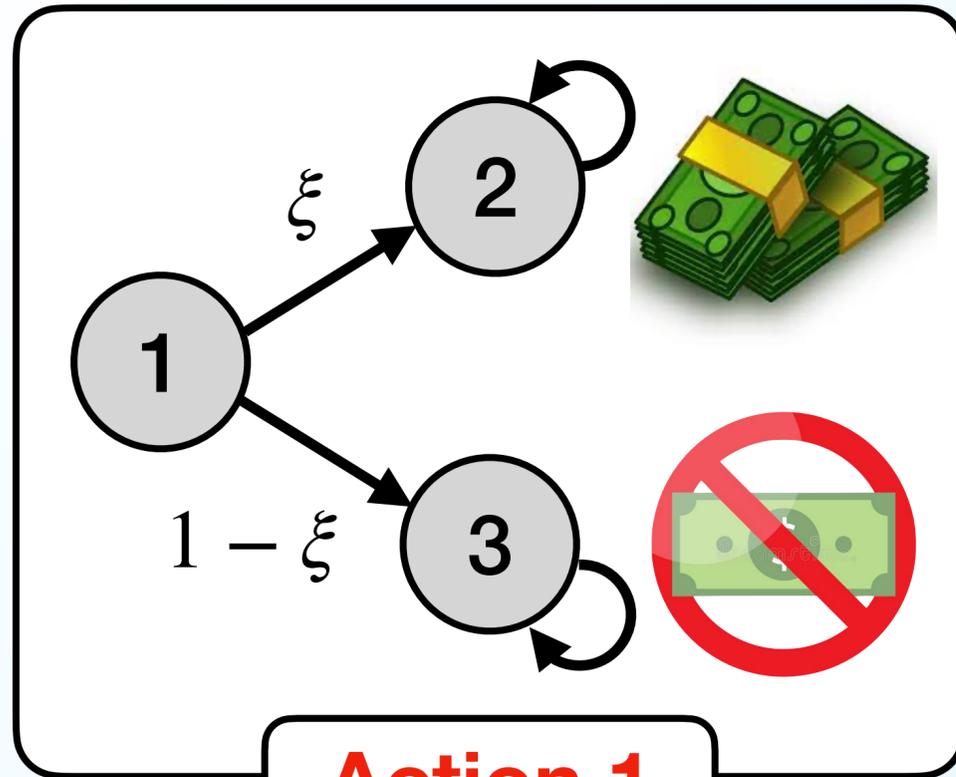


(s,a)-rectangular ambiguity sets

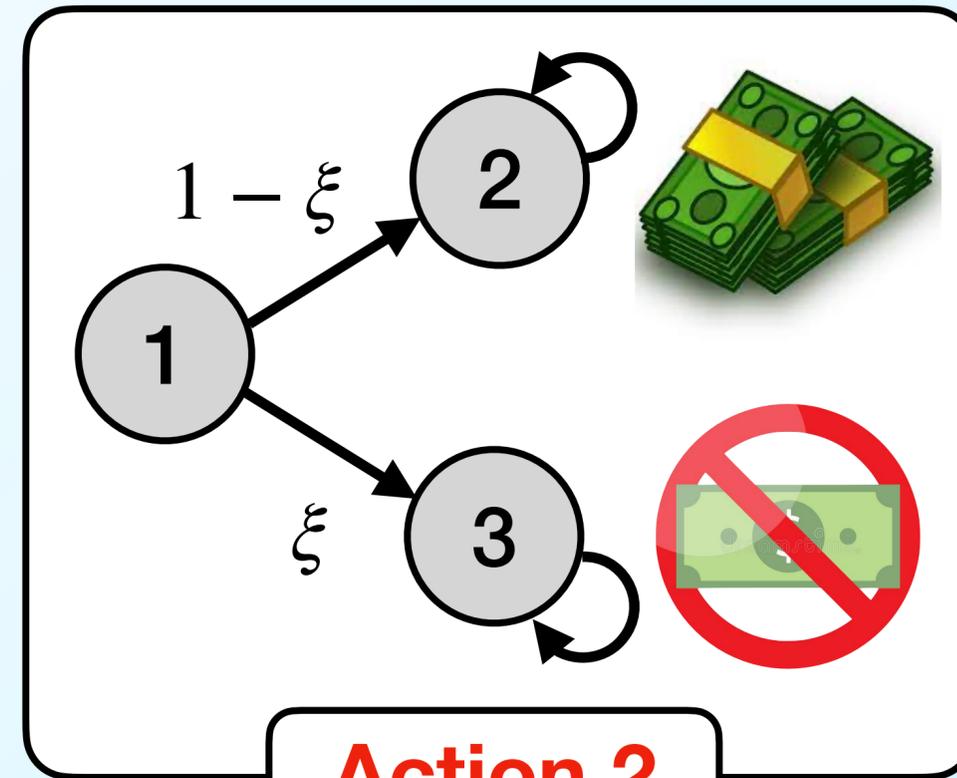
- 👍 Optimal policy **stationary** and **deterministic**
- 👍 Bellman optimality principle **holds**

General (non-rectangular) ambiguity sets

Example



Action 1



Action 2

for some unknown $\xi \in [0,1]$



Bellman optimality principle holds

s -Rectangular Ambiguity Sets: Bellman Operator

s-Rectangular Ambiguity Sets: Bellman Operator

Classical (non-robust) Bellman equations

$$v^*(s) = \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman equations

$$v^*(s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v^*(s') \right] \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

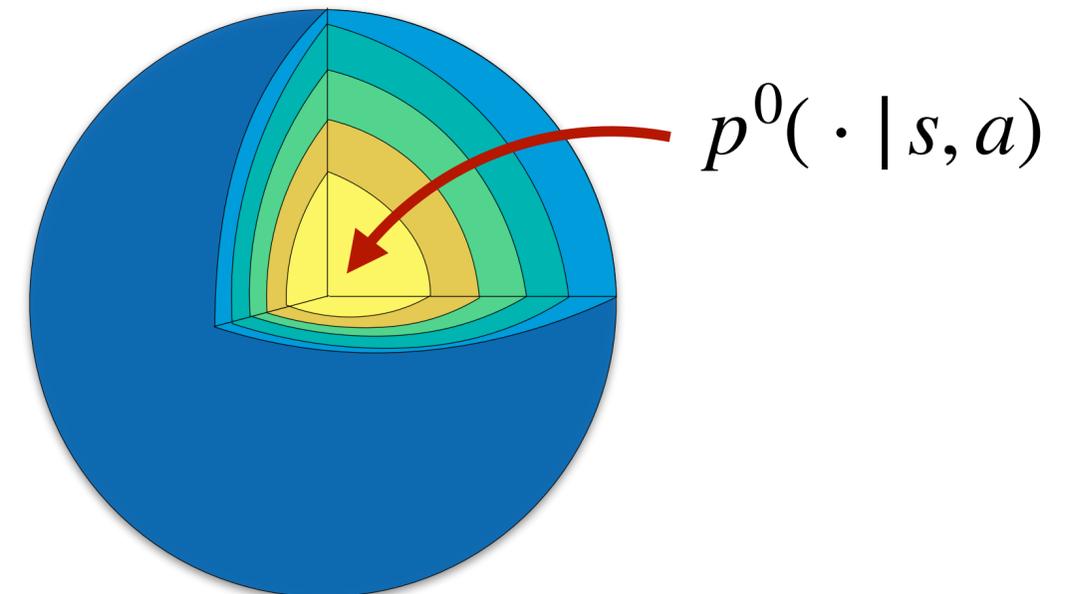
s-Rectangular Ambiguity Sets: Bellman Operator

Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$



s-Rectangular Ambiguity Sets: Bellman Operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\begin{aligned} [\mathfrak{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\ &\downarrow \\ &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \end{aligned}$$

Minimax theorem: exchange order of min and max

s-Rectangular Ambiguity Sets: Bellman Operator

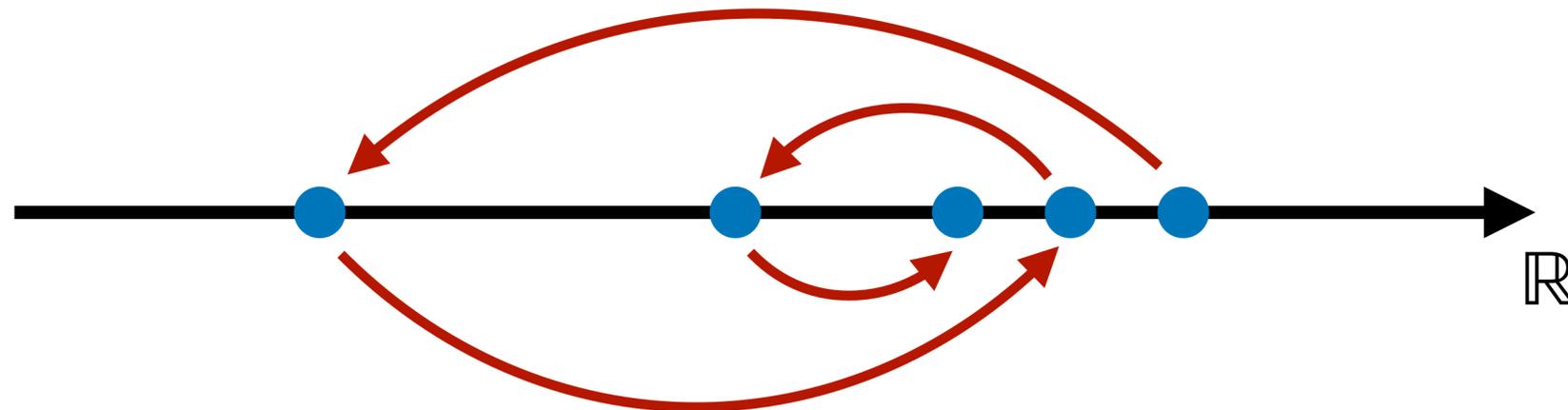
$$\begin{aligned} [\mathfrak{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\ &\downarrow \\ &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\ &\downarrow \\ &= \min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \end{aligned}$$

Linearity: we only need to consider ext $\Delta(\mathcal{A}) = \mathcal{A}$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\begin{aligned}
 [\mathfrak{B}v](s) &= \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\downarrow \\
 &= \min_{p \in \mathcal{P}_s} \max_{\pi \in \Delta(\mathcal{A})} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\} \\
 &\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?
 \end{aligned}$$

Bisection search:



s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\min_{p \in [\Delta(\mathcal{S})]^{\mathcal{A}}} \left\{ \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} : \sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\} \leq \theta$$

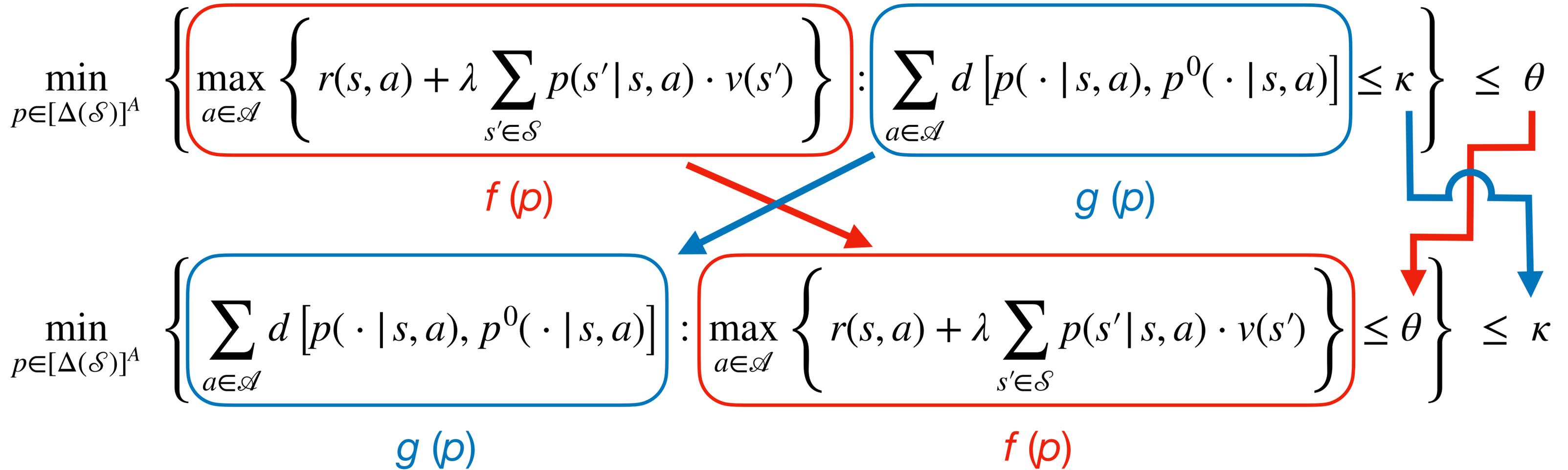
s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\min_{p \in [\Delta(\mathcal{S})]^{\mathcal{A}}} \left\{ \underbrace{\max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\}}_{f(p)} : \underbrace{\sum_{a \in \mathcal{A}} d[p(\cdot | s, a), p^0(\cdot | s, a)]}_{g(p)} \leq \kappa \right\} \leq \theta$$

s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$



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$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

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s-Rectangular Ambiguity Sets: Bellman Operator

$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

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$$\iff \sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d [p(\cdot | s, a), p^0(\cdot | s, a)] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$

Separability: of both objective and constraints in $a \in \mathcal{A}$

s-Rectangular Ambiguity Sets: Bellman Operator

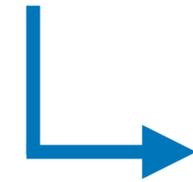
$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d \left[p(\cdot | s, a), p^0(\cdot | s, a) \right] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$

s-Rectangular Ambiguity Sets: Bellman Operator

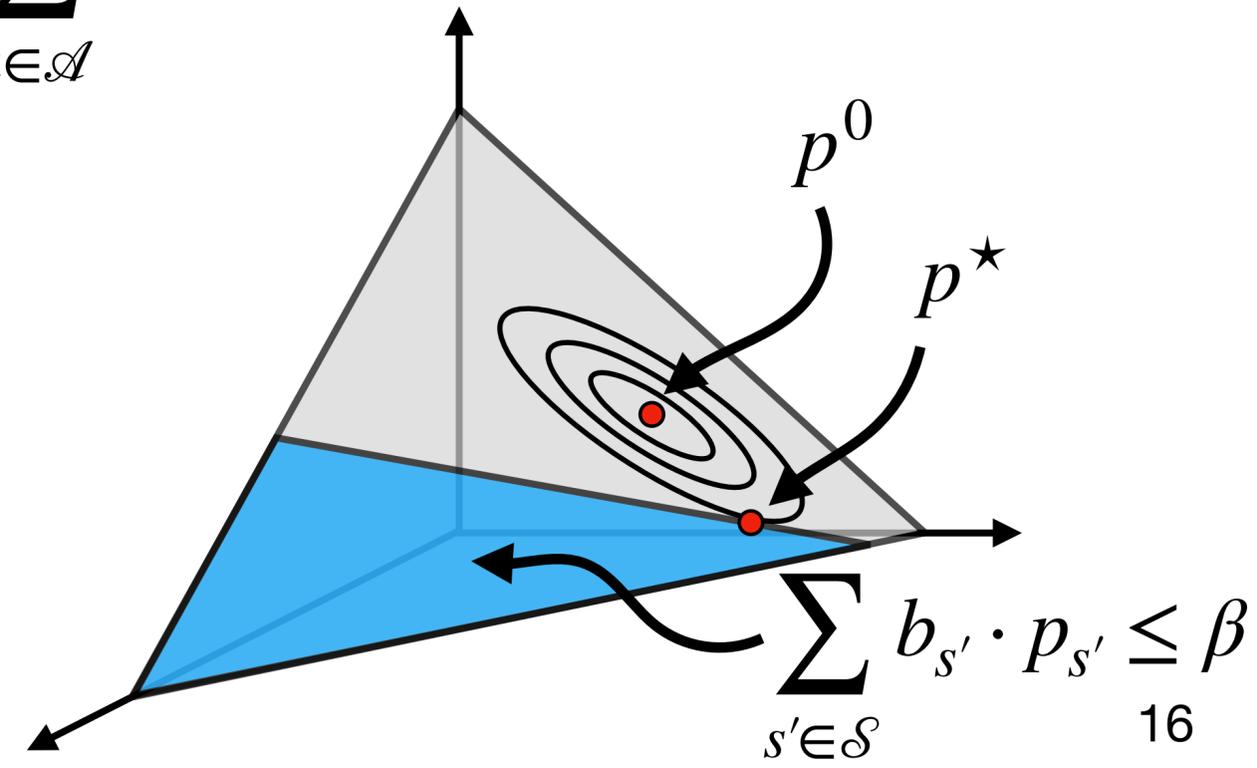
$$\min_{p \in \mathcal{P}_s} \max_{a \in \mathcal{A}} \left\{ r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right\} \leq \theta ?$$

$$\sum_{a \in \mathcal{A}} \min_{p_a \in \Delta(\mathcal{S})} \left\{ d \left[p(\cdot | s, a), p^0(\cdot | s, a) \right] : r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \leq \theta \right\} \leq \kappa$$



$$\Leftrightarrow \sum_{a \in \mathcal{A}} \mathfrak{P}(p^0; \lambda v, \theta - r(s | a)) \leq \kappa$$

with $\mathfrak{P}(p^0; b, \beta) =$ $\left[\begin{array}{l} \text{minimize} \\ p \\ \text{subject to} \\ d[p, p^0] \\ \sum_{s' \in \mathcal{S}} b_{s'} \cdot p_{s'} \leq \beta \\ p \in \Delta(\mathcal{S}) \end{array} \right]$



s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d [p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

s-Rectangular Ambiguity Sets: Bellman Operator

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Robust Bellman operator

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

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Projection problem

$$\mathfrak{P}(p^0; b, \beta) = \left[\begin{array}{l} \text{minimize} \\ p \\ \text{subject to} \end{array} \begin{array}{l} d [p, p^0] \\ \sum_{s' \in \mathcal{S}} b_{s'} \cdot p_{s'} \leq \beta \\ p \in \Delta(\mathcal{S}) \end{array} \right]$$

s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d [p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

Theorem

Assume \mathfrak{P} can be computed **exactly** in time $\mathcal{O}(h(S))$.
Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time
 $\mathcal{O}(AS \cdot h(S) \cdot \log[\bar{R}/\epsilon])$.

s-Rectangular Ambiguity Sets: Bellman Operator

Distance-constrained s-rectangular ambiguity set

$$\mathcal{P} = \prod_{s \in \mathcal{S}} \mathcal{P}_s \quad \text{with} \quad \mathcal{P}_s = \left\{ p(\cdot | s, \cdot) : \sum_{a \in \mathcal{A}} d [p(\cdot | s, a), p^0(\cdot | s, a)] \leq \kappa \right\}$$

Theorem

Assume \mathfrak{P} can be computed **exactly** in time $\mathcal{O}(h(S))$.

Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time $\mathcal{O}(AS \cdot h(S) \cdot \log[\bar{R}/\epsilon])$.

Assume \mathfrak{P} can be **computed to any accuracy** $\delta > 0$ in time $\mathcal{O}(h(\delta))$. Then \mathfrak{B} can be **computed to accuracy** $\epsilon > 0$ in time $\mathcal{O}(AS \cdot h(\epsilon\kappa/[2A\bar{R} + A\epsilon]) \cdot \log[\bar{R}/\epsilon])$.

s-Rectangular Ambiguity Sets: Bellman Operator

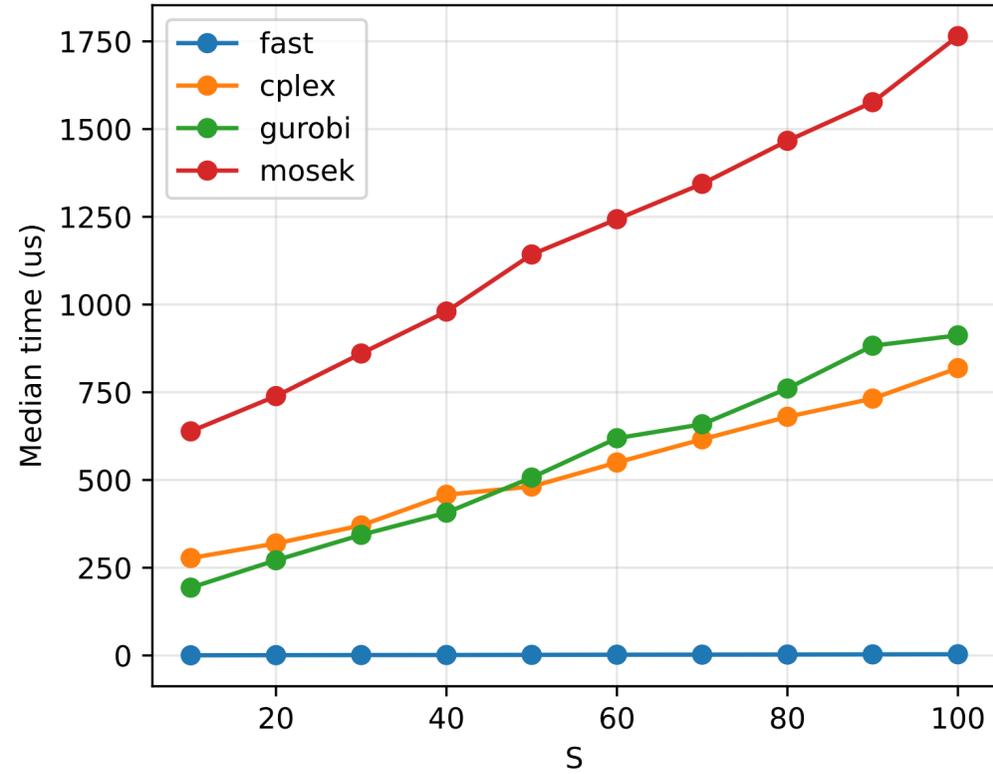
Divergence	$d_a(\cdot, p^0)$	Ours	Previous
KL-Divergence	$\sum_{s' \in \mathcal{S}} p(s' s, a) \cdot \log \left(\frac{p(s' s, a)}{p^0(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	$\mathcal{O}(\ell^2 \cdot S^2 \cdot A)$
Burg Entropy	$\sum_{s' \in \mathcal{S}} p^0(s' s, a) \cdot \log \left(\frac{p^0(s' s, a)}{p(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	(none)
Variation Distance	$\sum_{s' \in \mathcal{S}} p(s' s, a) - p^0(s' s, a) $	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^2 A \cdot \log S)$
χ^2 -Distance	$\sum_{s' \in \mathcal{S}} \frac{[p(s' s, a) - p^0(s' s, a)]^2}{p^0(s' s, a)}$	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^{4.5} \cdot A)$

s-Rectangular Ambiguity Sets: Bellman Operator

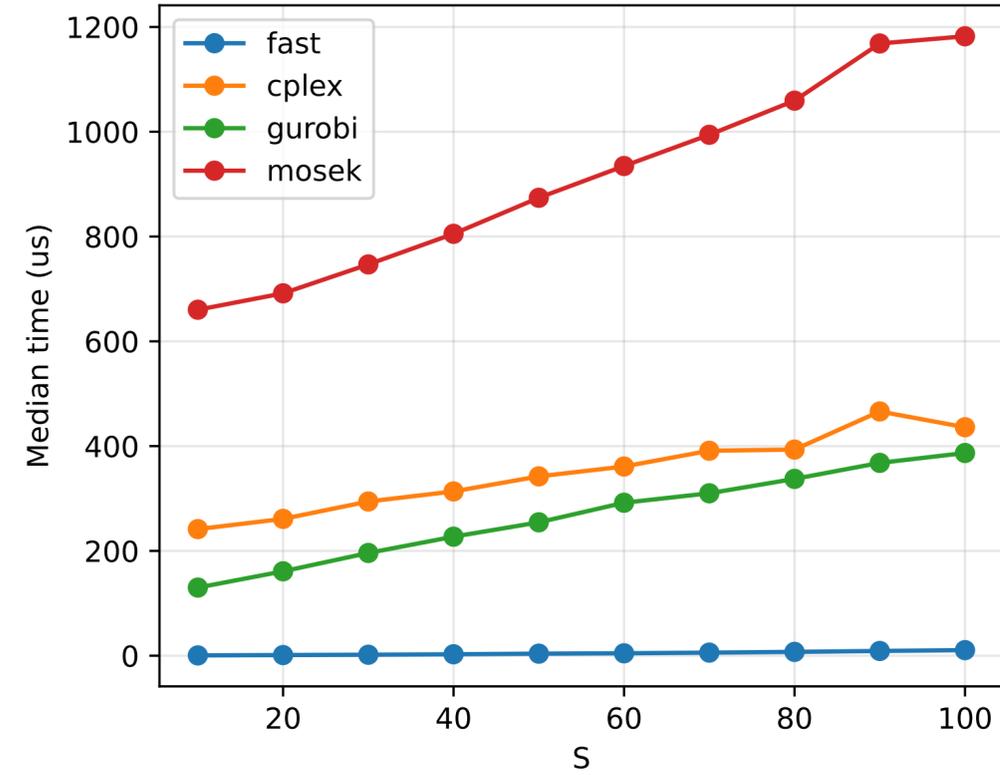
Divergence	$d_a(\cdot, p^0)$	Ours	Previous
KL-Divergence	$\sum_{s' \in \mathcal{S}} p(s' s, a) \cdot \log \left(\frac{p(s' s, a)}{p^0(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	$\mathcal{O}(\ell^2 \cdot S^2 \cdot A)$
Burg Entropy	$\sum_{s' \in \mathcal{S}} p^0(s' s, a) \cdot \log \left(\frac{p^0(s' s, a)}{p(s' s, a)} \right)$	$\mathcal{O}(S^2 A \cdot \log A)$	(none)
Variation Distance	$\sum_{s' \in \mathcal{S}} p(s' s, a) - p^0(s' s, a) $	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^2 A \cdot \log S)$
χ^2 -Distance	$\sum_{s' \in \mathcal{S}} \frac{[p(s' s, a) - p^0(s' s, a)]^2}{p^0(s' s, a)}$	$\mathcal{O}(S^2 A \cdot \log S)$	$\mathcal{O}(S^{4.5} \cdot A)$

s-Rectangular Ambiguity Sets: Projection Problem

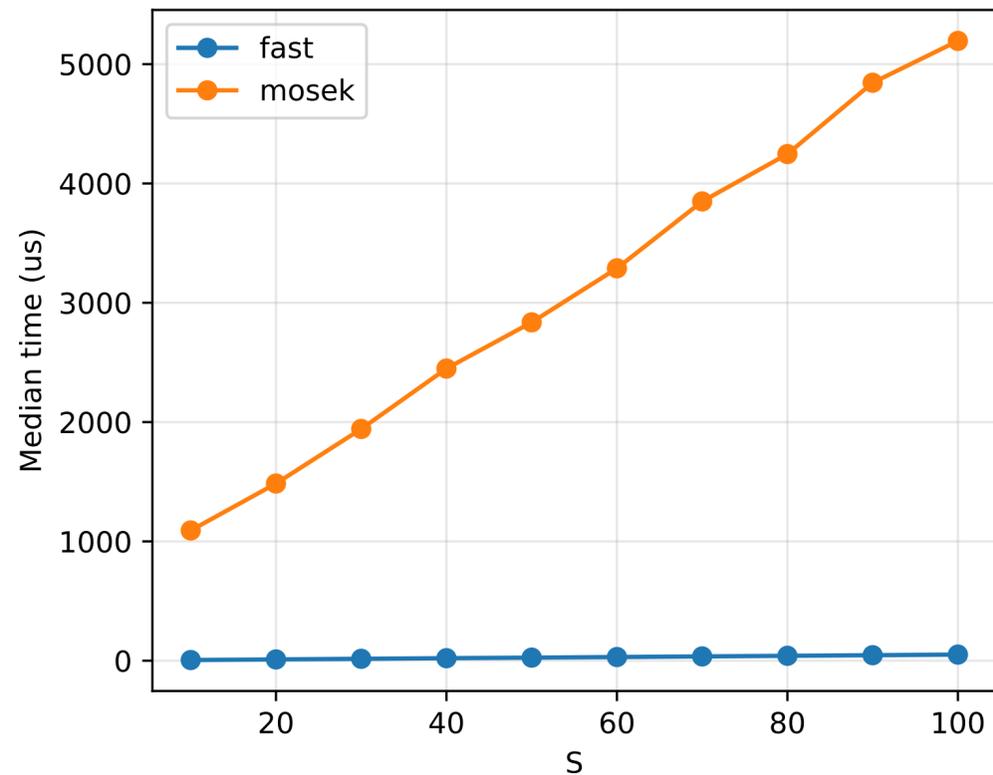
1-norm



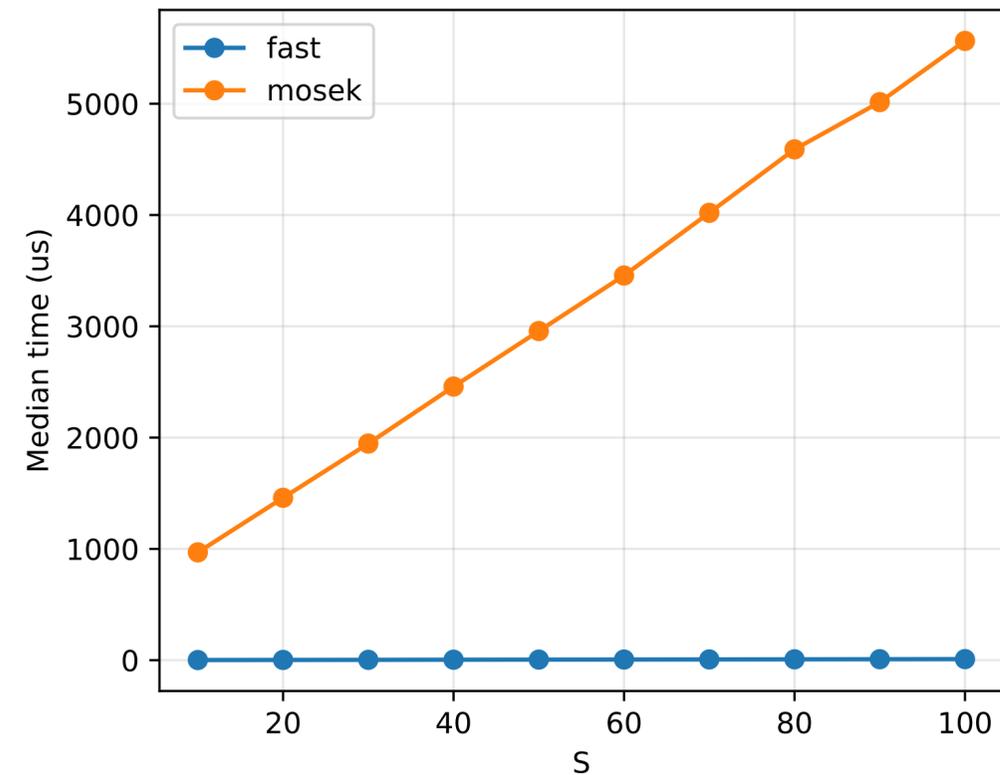
2-norm



KL-Div

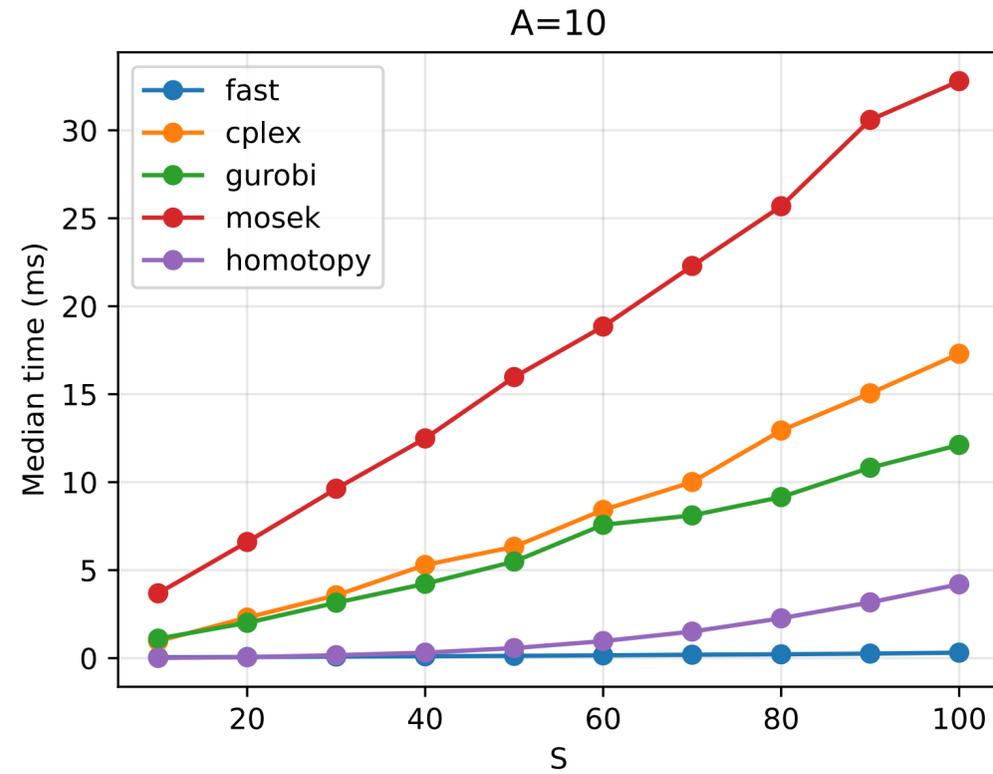


Burg

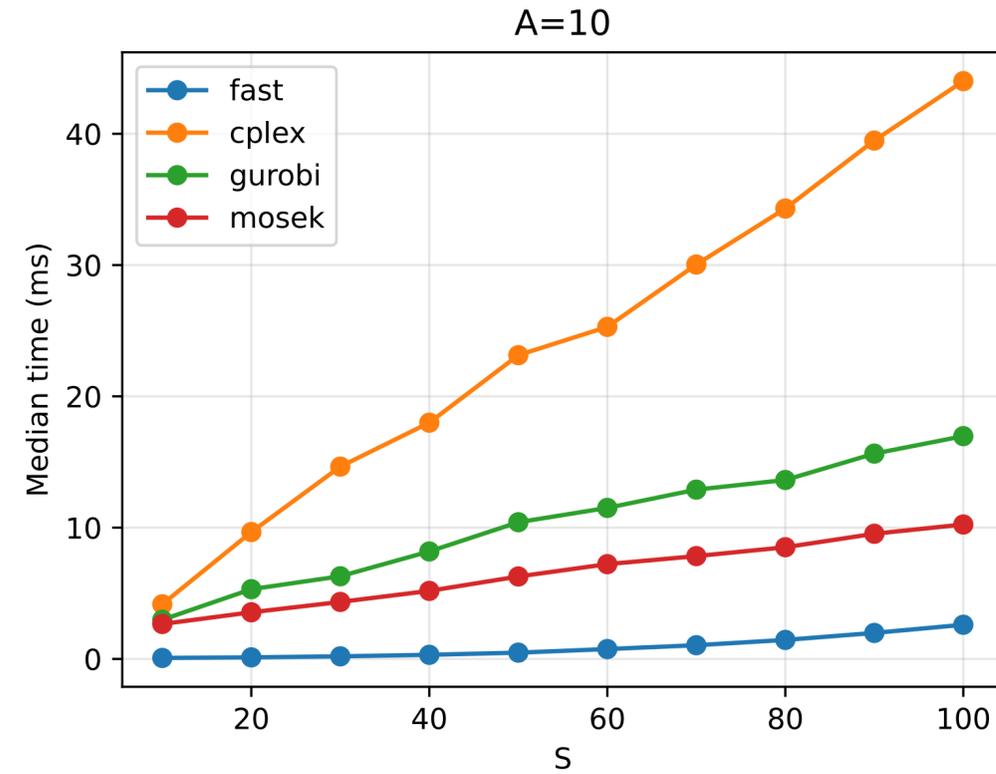


s-Rectangular Ambiguity Sets: Bellman Operator

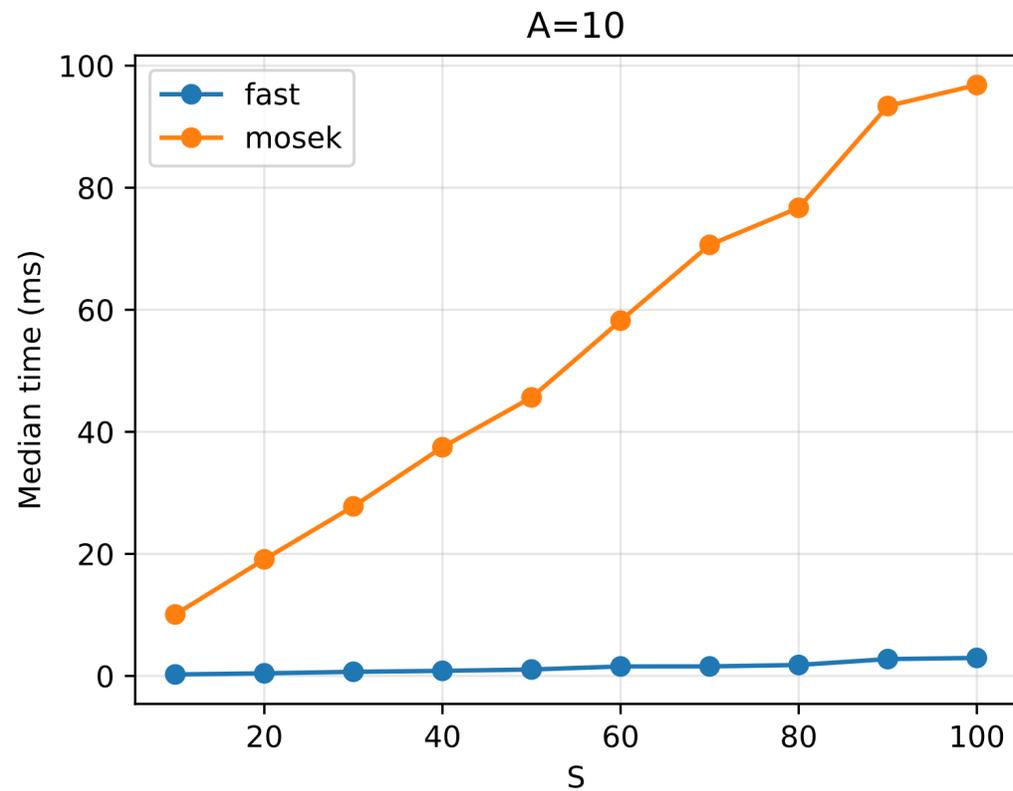
1-norm



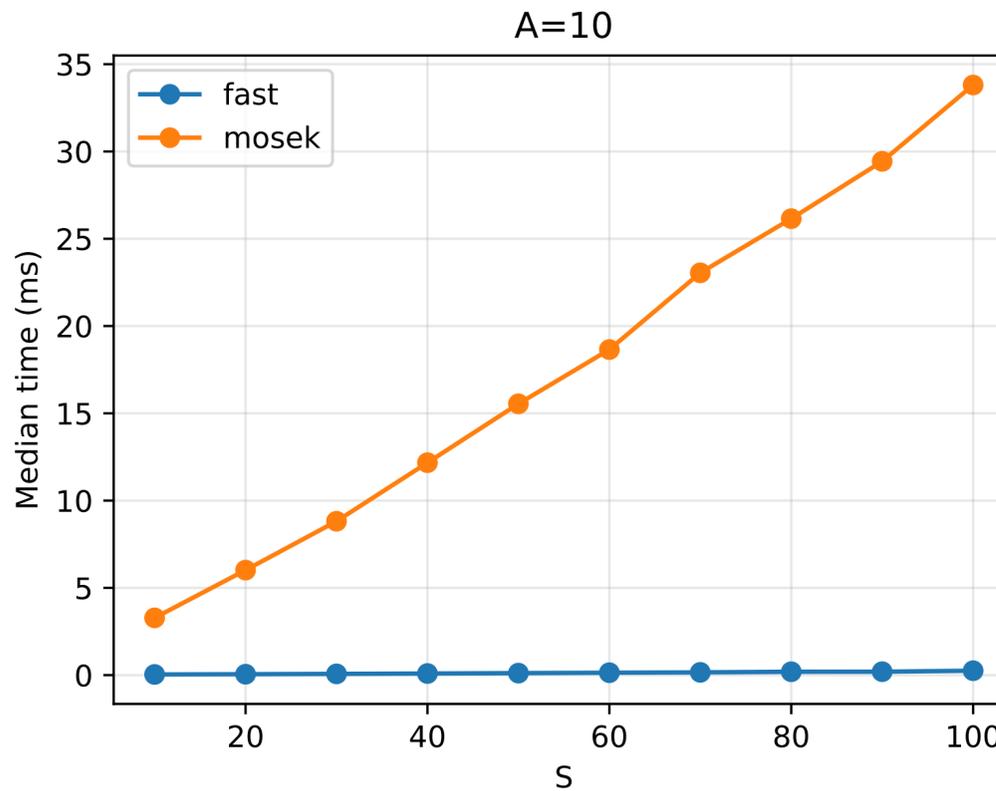
2-norm



KL-Div

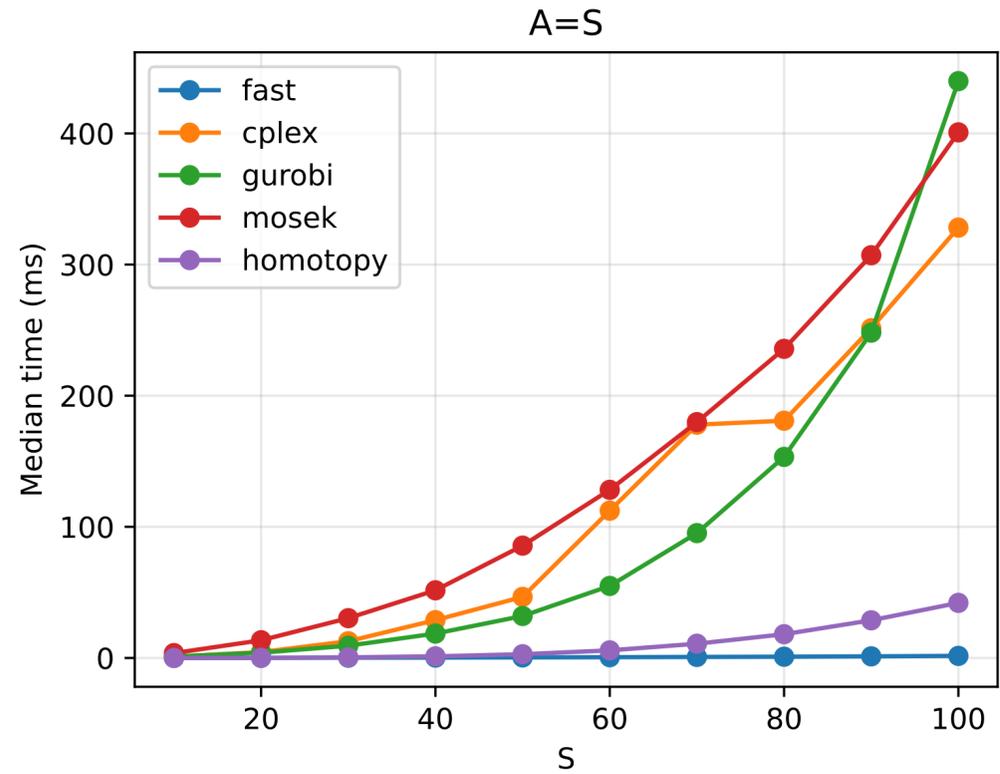


Burg

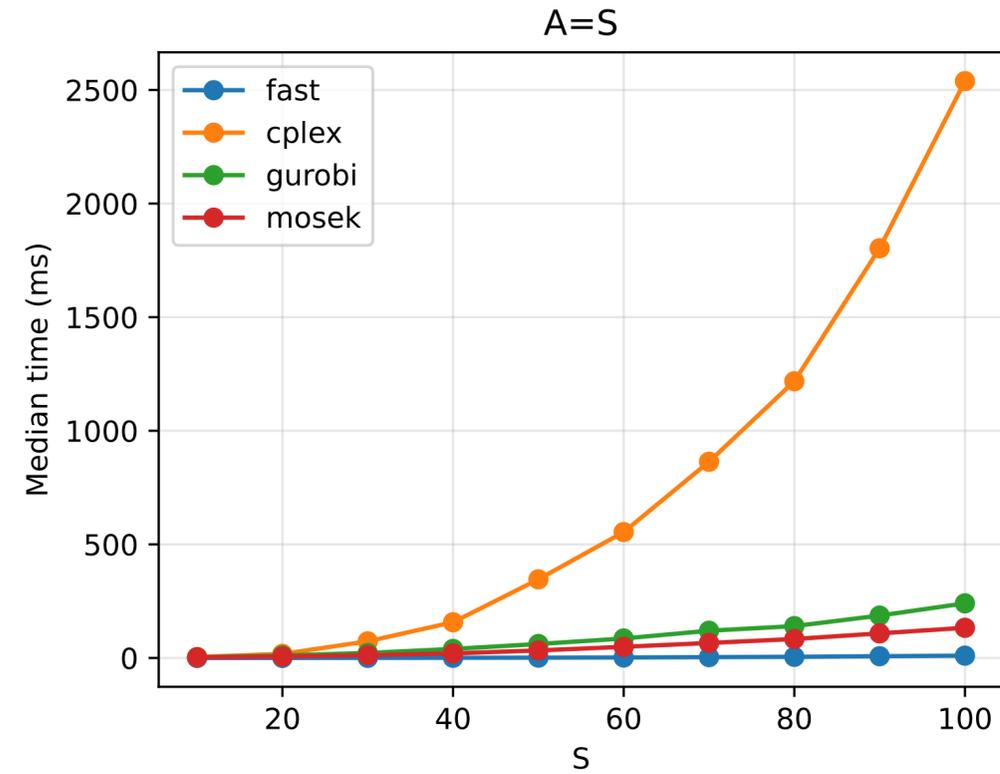


s-Rectangular Ambiguity Sets: Bellman Operator

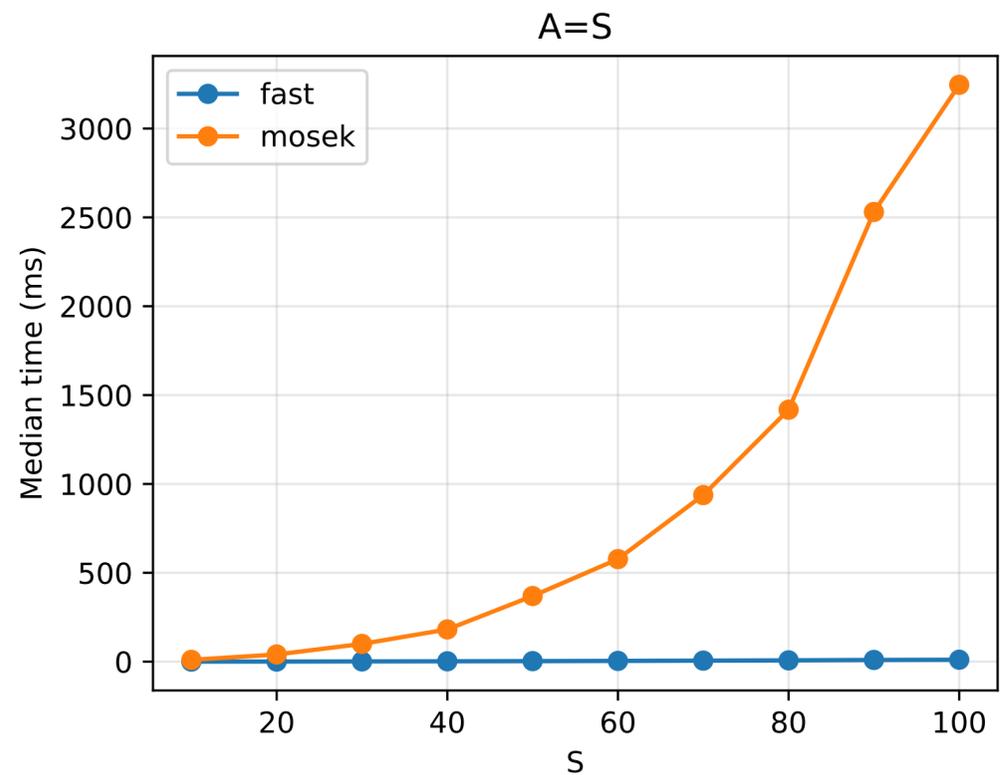
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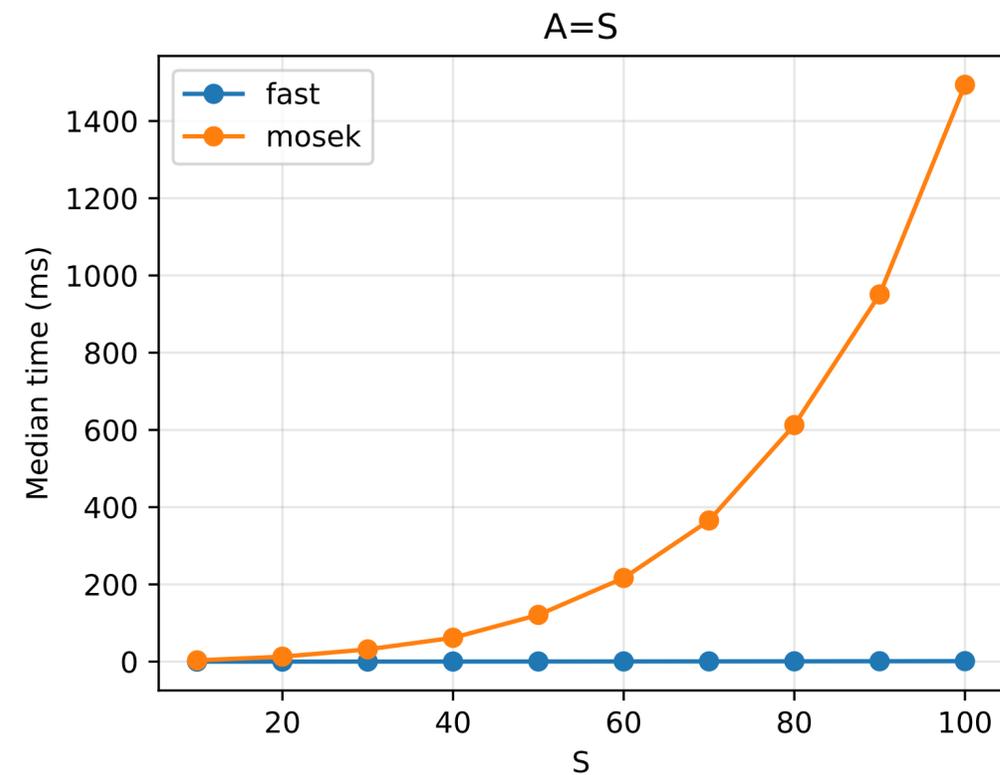
2-norm



KL-Div



Burg



s-Rectangular Ambiguity Sets: Partial Policy Iteration

Reconsider idea of **(modified) policy iteration**:

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Reconsider idea of **(modified) policy iteration**:

Policy improvement

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$



expensive operator: requires **robust Bellman operator**

s-Rectangular Ambiguity Sets: Partial Policy Iteration

Reconsider idea of **(modified) policy iteration**:

Policy improvement

$$[\mathfrak{B}v](s) = \max_{\pi \in \Delta(\mathcal{A})} \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

 **expensive operator**: requires **robust Bellman operator**

Policy evaluation

$$[\mathfrak{B}(\pi)v](s) = \min_{p \in \mathcal{P}_s} \left\{ \sum_{a \in \mathcal{A}} \pi(a) \cdot \left[r(s, a) + \lambda \sum_{s' \in \mathcal{S}} p(s' | s, a) \cdot v(s') \right] \right\}$$

 **cheaper operator**: can be recast as Bellman operator of a nominal MDP

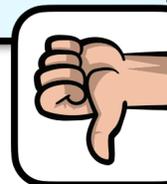
s-Rectangular Ambiguity Sets: Robust Modified Policy Iteration

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Repeat: starting with $i = 0$, v^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathfrak{B}v^i$ and let π^{i+1} be the corresponding greedy policy



expensive operator:
robust value iteration

$i = i + 1$

Until $\|w^{i+1} - \mathcal{Q}^{i+1,N}\|_{\infty} < \frac{1 - \lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: Robust Modified Policy Iteration

Repeat: starting with $i = 0$, v^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathfrak{B}v^i$ and let π^{i+1} be the corresponding greedy policy

Policy evaluation

Compute sequence $q^{i+1,j+1} = \mathfrak{B}(\pi^{i+1})q^{i+1,j}$ with $q^{i+1,0} = w^{i+1}$ until $\|q^{i+1,j+1} - q^{i+1,j}\|_\infty \leq (1 - \lambda)\epsilon_{i+1}$

$i = i + 1$



cheaper operator:
no maximum involved

Until $\|w^{i+1} - q^{i+1,N}\|_\infty < \frac{1 - \lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: **Partial Policy Iteration**

Repeat: starting with $i = 0$, v^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathfrak{B}v^i$ and let π^{i+1} be the corresponding greedy policy

Policy evaluation

Alternate between **single** robust Bellman evaluation $\mathfrak{B}(\pi^{i+1})$ and **multiple** nominal Bellman evaluations under **worst-case** p .

$i = i + 1$



cheap (!) operator:
nominal evaluations

Until $\|w^{i+1} - \mathfrak{Q}^{i+1,N}\|_{\infty} < \frac{1 - \lambda}{2} \cdot \delta$

s-Rectangular Ambiguity Sets: **Partial Policy Iteration**

Repeat: starting with $i = 0$, v^0 arbitrary

Policy improvement

Compute $w^{i+1} = \mathfrak{B}v^i$ and π^{i+1} the corresponding greedy policy

Theorem

Assume $\epsilon_{i+1} < \lambda^c \cdot \epsilon_i$ for some $c > 1$. Then the optimality gap of partial policy iteration satisfies:

$$\|v(\pi^{i+1}) - v^*\|_\infty \leq \lambda^i \left(\|v(\pi^1) - v^*\|_\infty + \frac{2\epsilon_1}{(1 - \lambda^{c-1})(1 - \lambda)} \right)$$

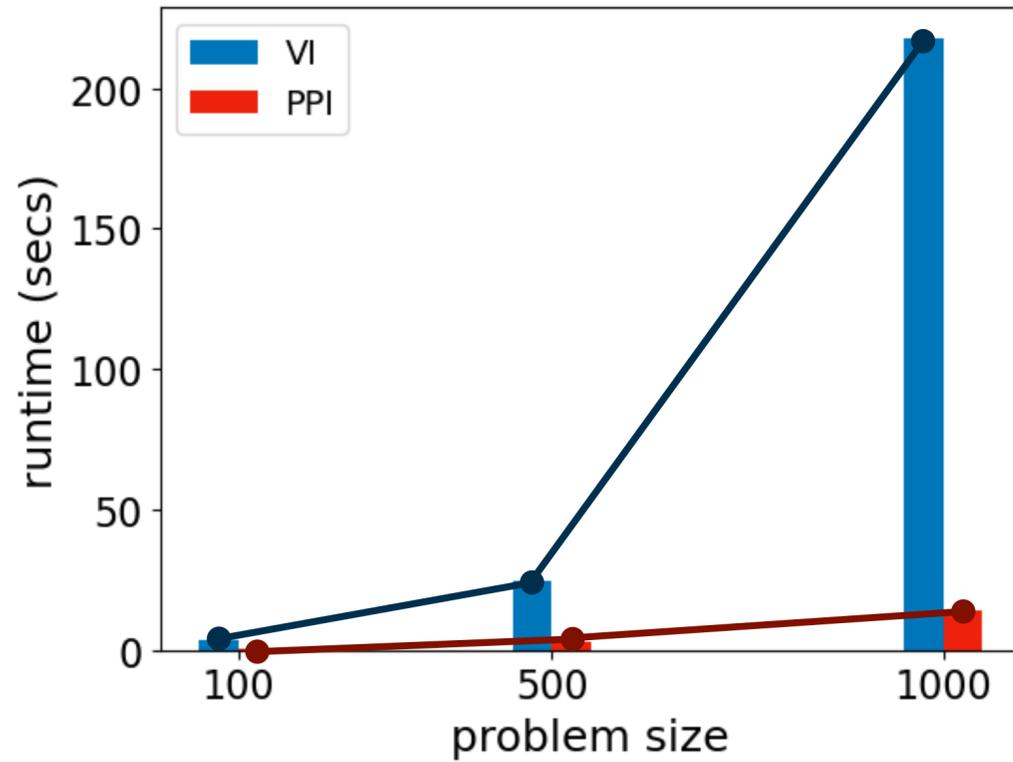
$i = i + 1$

Until $\|w^{i+1} - \mathfrak{D}^{i+1,N}\|_\infty < \frac{1 - \lambda}{2} \cdot \delta$

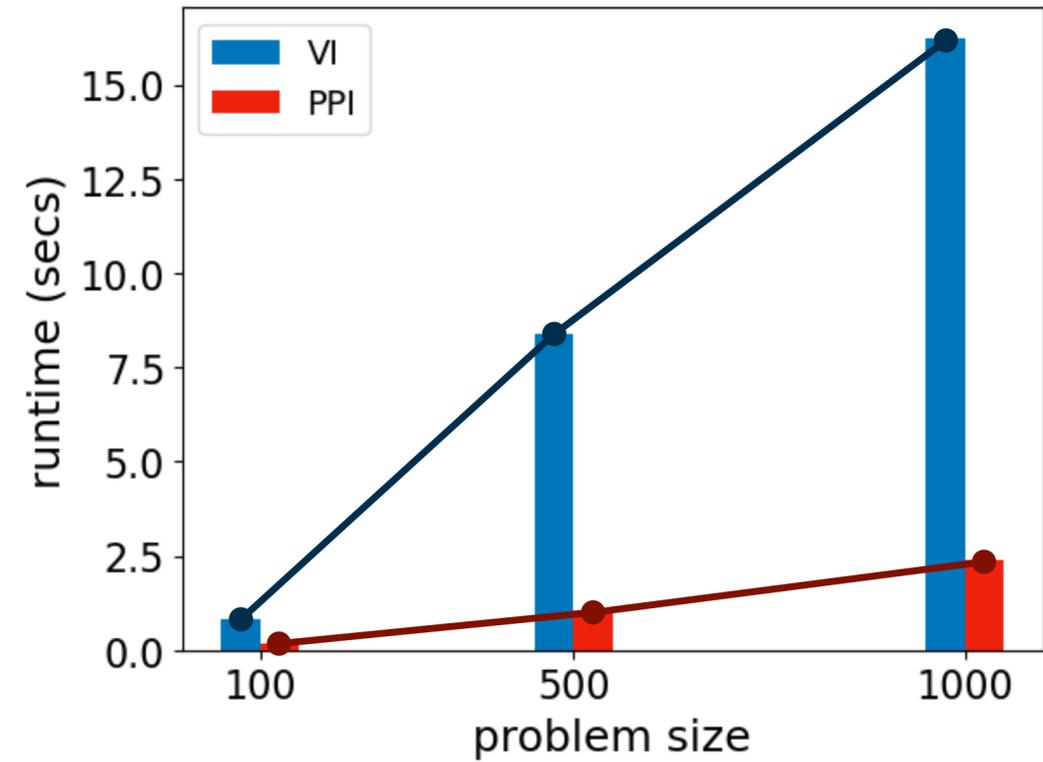
s-Rectangular Ambiguity Sets: **Partial Policy Iteration**

Inventory

uniform

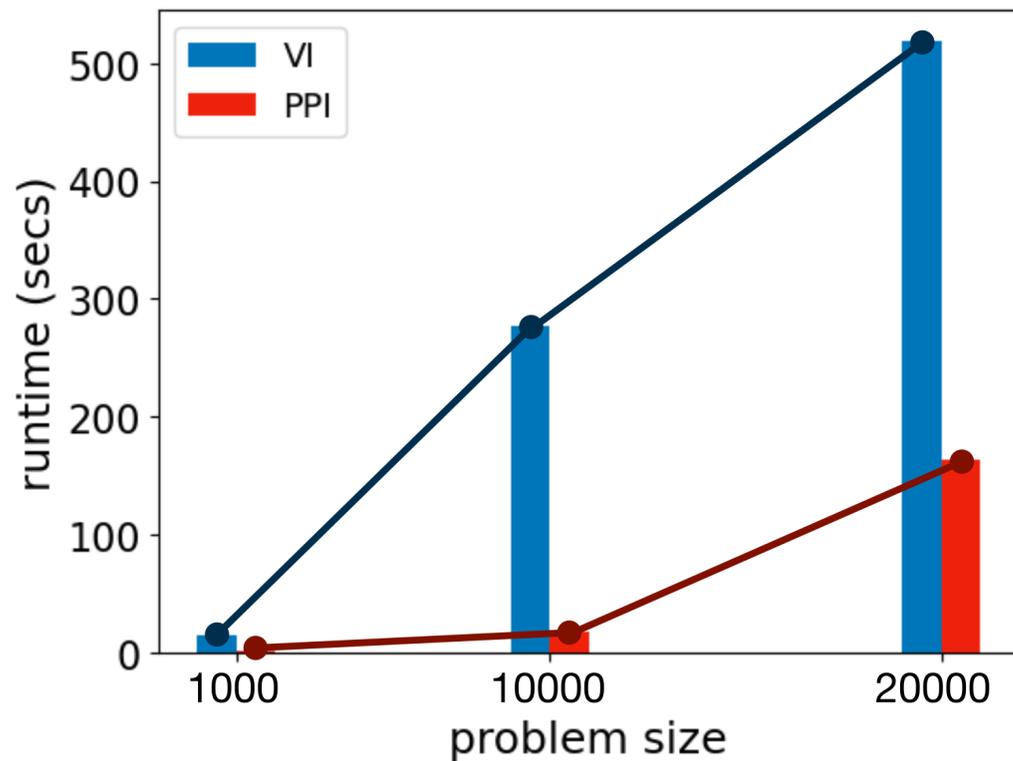


weighted

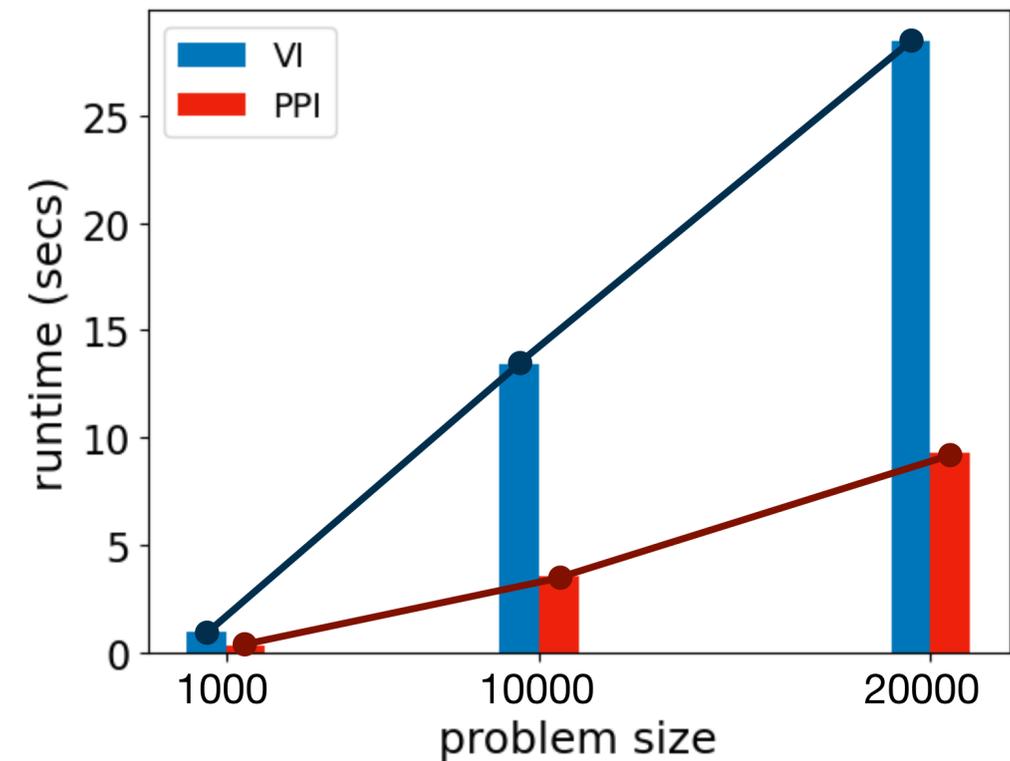


Cart Pole

uniform



weighted



Conclusions: **MDPs** – **Now More Important Than Ever!**

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 **ChatGPT** ✨ **Gemini** 🌟 **Claude**

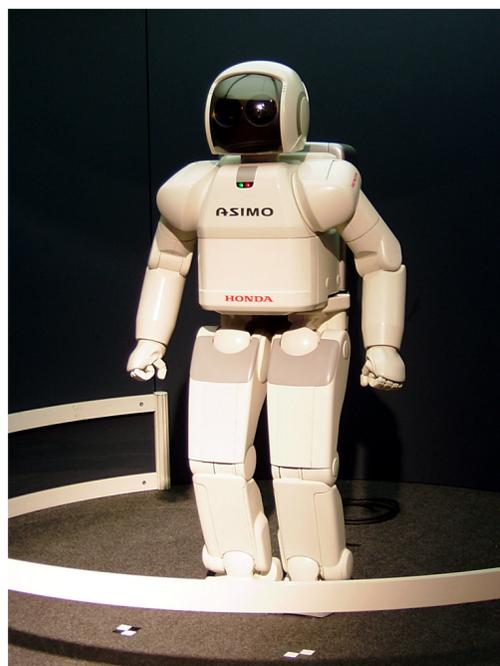
Reinforcement learning with human feedback

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Reinforcement learning with human feedback



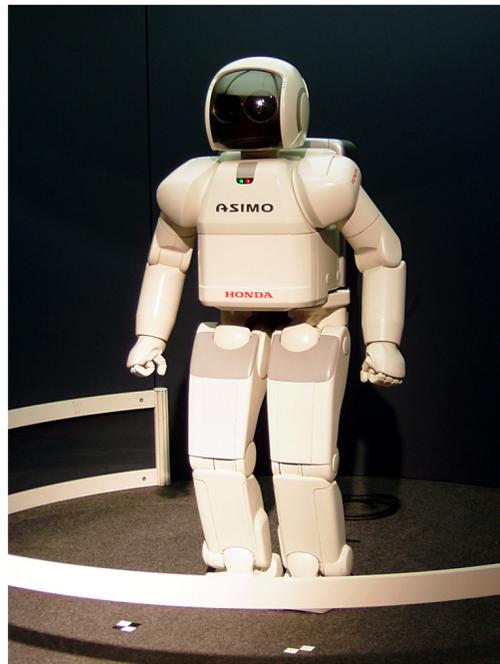
Robotics

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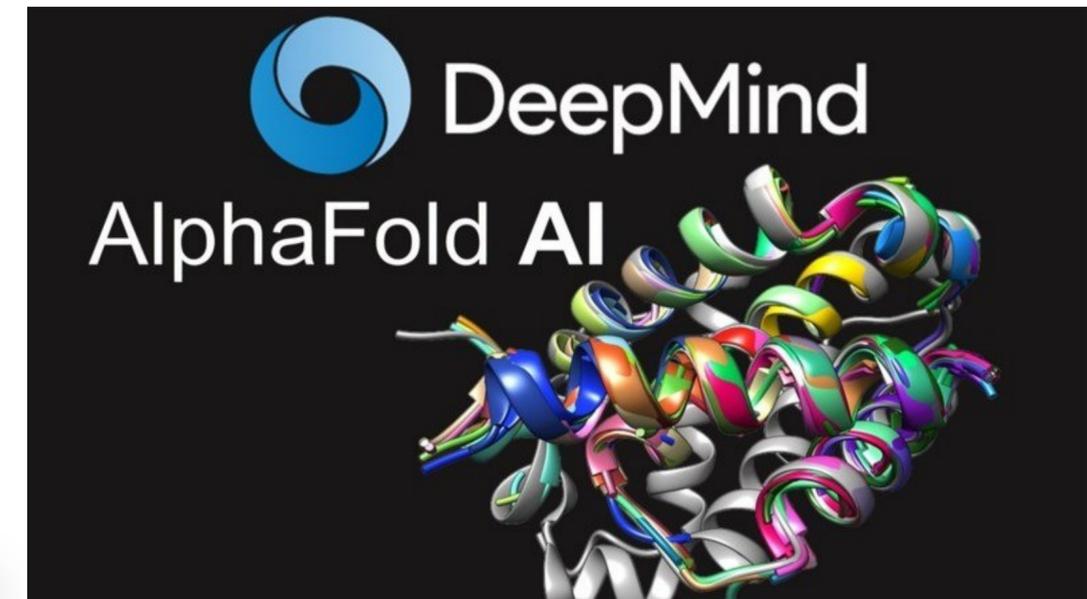
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Reinforcement learning with human feedback



Robotics



Medical discovery

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- **MDPs** give us the **mathematical understanding** that underpins all this.
- There are **exciting new developments** that **bring both fields closer**:
 - **Linear MDPs**: Reinforcement learning with linear dynamics
 - **Weakly Coupled MDPs, Factored MDPs, Constrained MDPs, Safe RL, ...**
 - covered extensively in **top AI/ML conferences** in recent years!

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